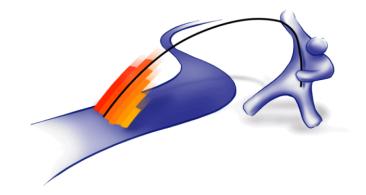


Carrier Grade NAT Deployment, Traffic Impact, and Logging Considerations



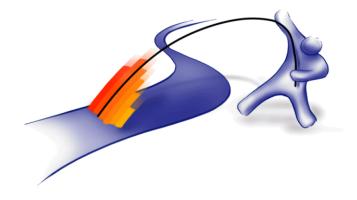
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13.03.2012



Agenda

- Motivation
- NAT overview
- Implementation and Deployment
- NAT Logging
- Traffic impact
- Summary



Motivation

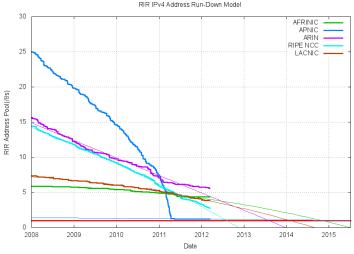
Problem

Increasing number of IP-enabled (mobile) devices

→ IPv4 Address space exhausted

Solutions

- 1. Efficient use of remaining IPv4 addresses
- 2. Migration to IPv6 (+translations to reach IPv4)
- → **both** require address translation!



http://www.potaroo.net/tools/ipv4/

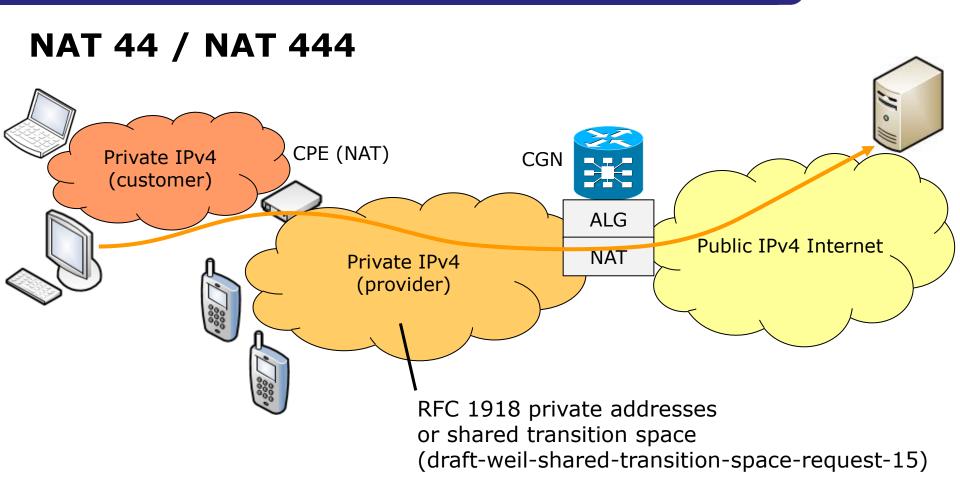
Implementation

- Address translation in the provider network
- Carrier Grade Network Address Translation (CGN), also Large Scale NAT (LSN)

...CGNs will exist for several years (forever?)

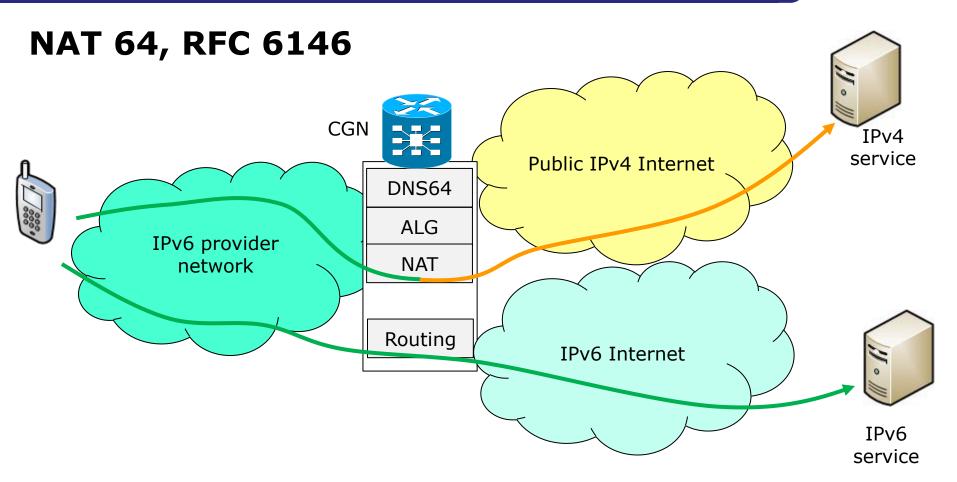
...should use IPv4 addresses efficiently (many subscribers per public IPv4 Address)

NAT overview



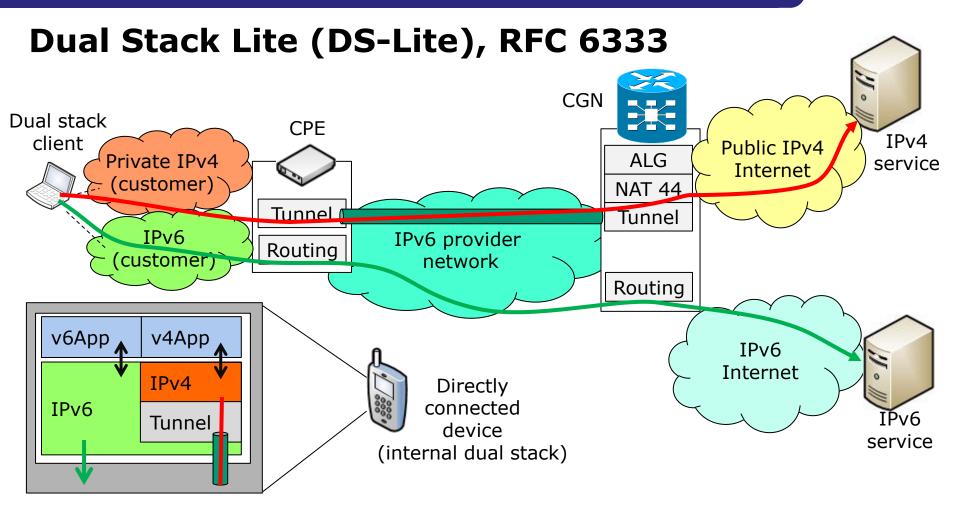
- More efficient usage of IPv4 resources
- Application Layer Gateway (ALG) for IP-address-bound applications
- Short-term solution no IPv6 deployed

NAT overview



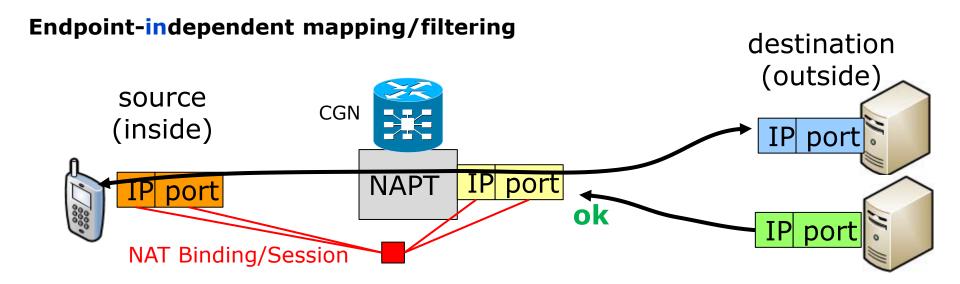
- Provider network fully migrated to IPv6
- IPv6 clients only
- IPv4 content reachable via DNS64-translation and NAT
- Long-term solution will be there for several years (?)

NAT overview



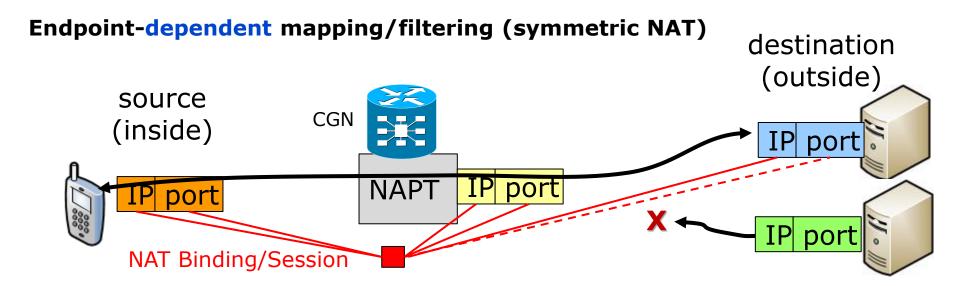
- Dual Stack from client point of view
 - IPv4 and/or IPv6 clients/applications
- Provider network IPv6 only

Network Address and Port Translation (NAPT) Implementations, RFC 4787



- Any endpoint can access client through NAT binding
- Prerequisite for NAT-traversal mechanisms (STUN, P2P applic., ...)
- RFC 4787: MUST use endpoint-independent mapping

Network Address and Port Translation (NAPT) Implementations, RFC 4787



- Only initially addressed endpoint can access client through NAT binding
- Breaks NAT-traversal mechanisms (STUN, P2P, ...) ... use it deliberately?
- Current discussions: higher efficiency of endpoint-dependent mapping

Operational Requirements

NAT Efficiency

- Efficient resource (IPv4 address) usage
- Free used resources as soon as possible
- → **Timeouts** required

Security

- Several Address/port scans could exhaust public IP address pool easily
- → Limit maximum resource usage: **portlimit**

Availability

→ **Failover** capabilities

Logging

- How to track down malicious activities?
- Public IP Address is no longer a "temporary user identifier"
- → Private IP ↔ public IP mapping required in logs

Exemplary CGN device: Cisco CGSE for CRS-1

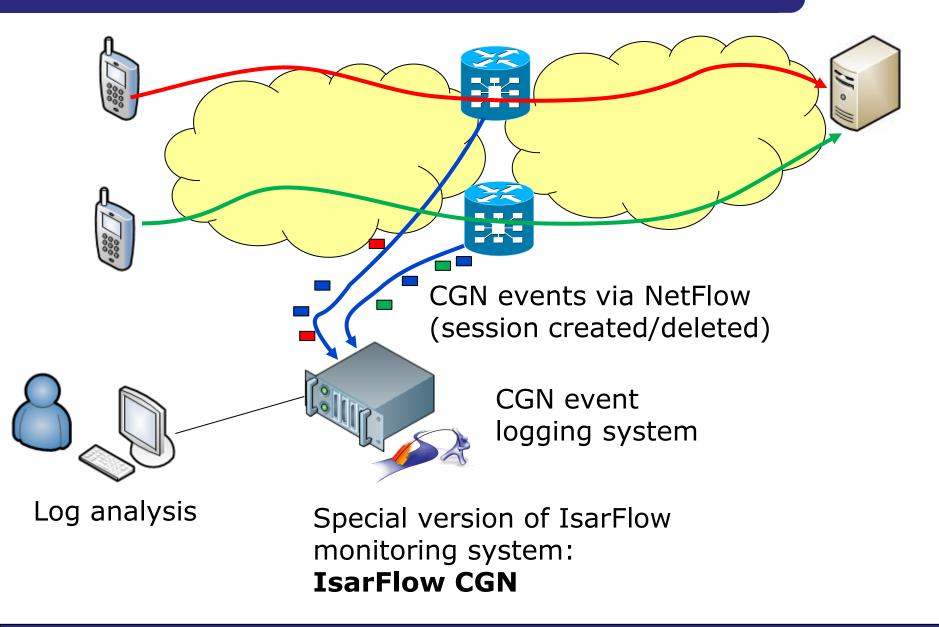
Performance

- Concurrent sessions: 20 Million
- Session creation rate: 1 Msessions/s
- Throughput: 20 Gbps full-duplex

Default configuration

- Portlimit 100
- Timeouts (inital/active)
 - TCP timeout 120s/1800s
 - UDP timeout 30s/120s
 - ICMP timeout: 60 s





Logging requirements and solutions

Worst case scenario for dimensioning: failover

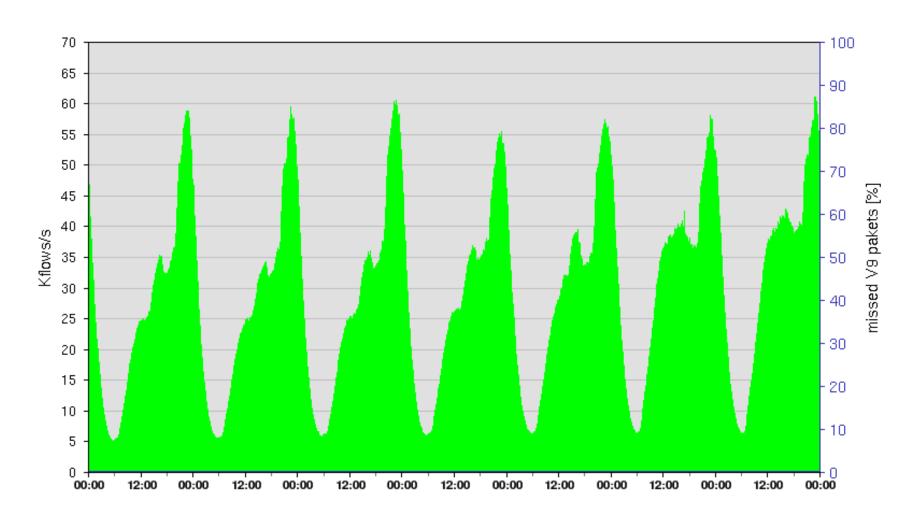
- All sessions of one location move to other locations
- Event-burst from one CGSE
 - 20 Million add events in 20 seconds
 - 1 Mevents/s → **180 Mbit/s**



Performance of IsarFlow CGN

- NAT 44 performance per server (COTS Linux box)
 - 1.5 Mevents/s → **270 Mbit/s** max. sustained rate
 - Loading of data into compressed database at that rate
 - Peak rate beyond 3 Mevents/s (540 Mbit/s) without loss
- Performance scales with number of servers (distributed DB)
- Storage requirements: 8 MB for 1 Million sessions

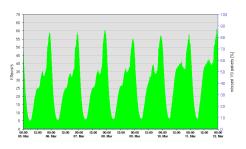
Typical event rate patterns



Ideas for reducing logging effort

Bulk port allocation

- Allocate several ports at once for each client
- → One log event for large port range



Problem: deterministic source ports are a security problem (RFC 6056)

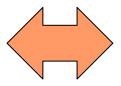
Possible solution: algorithmic port scattering

General problem of bulk port allocation

Bulk allocation is "port over-provisioning" → lower NAT efficiency

Trade-off

NAT efficiency (cost of public IPv4 addresses)



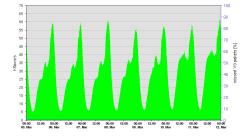
Logging efficiency

(cost of hard disk space)

Traffic impact

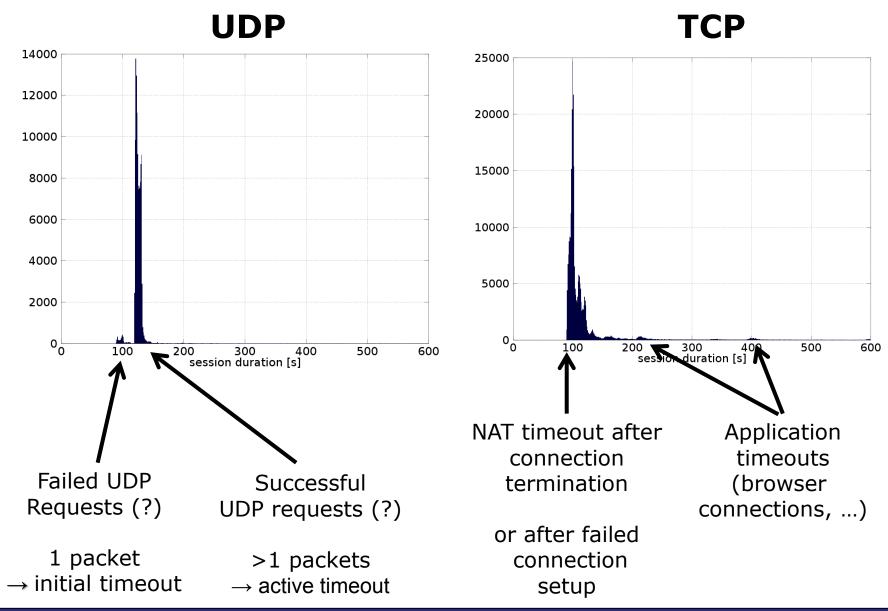
First studies of traffic characteristics

Based on inspection of 400k sessions



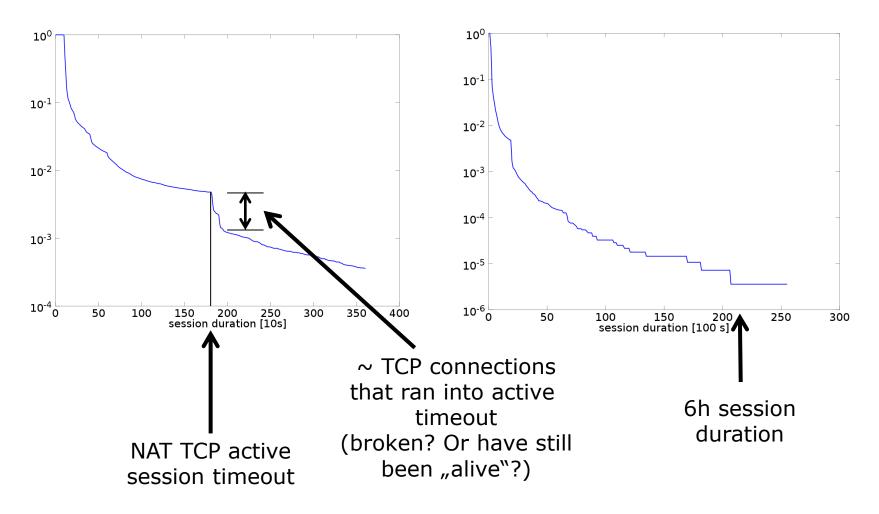
- No surprise: lots of short sessions
 - → Cause high NAT event rates
 - → Affect NAT efficiency: timeout until port can be reused
- Avg. Session duration (incl. timeouts) 135 s
- > 30 % UDP sessions (!)
 - DNS Resolver in public IP space?
 - SIP/VoIP also across NAT?

Traffic impact



Traffic impact

TCP - CCDF



Summary & Outlook

Summary

- Carrier grade NATs are currently getting deployed
- Still lots of standardization effort and new ideas (IETF Softwire, Behave, v6ops...)
- NAT session logging is a major concern
 - IsarFlow CGN proves feasibility of large-scale full session logging
 - Bulk allocation: Trade-off between NAT efficiency and logging
- Traffic impact
 - Portlimit: limited number of sessions per user
 - Timeouts: keepalives necessary, as with current CPE NATs
 - Short sessions: Cause high event rates and block resources due to timeouts

Outlook

- How will mobile session behavior evolve?
 - Impact on NAT efficiency
 - Impact on NAT configuration (timeouts, ...)
 - Impact on Logging requirements



References

IETF documents

- RFC 4787: Network Address Translation (NAT) Behavioral Requirements for Unicast UDP
- draft-ietf-behave-Isn-requirements
- RFC 6146: Stateful NAT64: Network Address and Protocol Translation from IPv6 Clients to IPv4 Servers
- RFC 6535: Dual-Stack Hosts Using "Bump-in-the-Host" (BIH)
- RFC 6144: Framework for IPv4/IPv6 Translation
- RFC 6052: IPv6 Addressing of IPv4/IPv6 Translators
- RFC 5382: NAT Behavioral Requirements for TCP
- RFC 6056: Recommendations for Transport-Protocol Port Randomization
- RFC 6333: Dual-Stack Lite Broadband Deployments Following IPv4 Exhaustion
- RFC 6269: Issues with IP Address Sharing