

Design and Evaluation of Scheduling Algorithms for LTE Femtocells

(Entwurf und Bewertung von Ressourcenzuweisungen für LTE Femtozellen)

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Diplomarbeit von Kay Henzel

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Overview

- Introduction and Motivation
- Radio Resource Management Problem
- Radio Resource Assignment Algorithms
- Simulation Scenario and Results
- Summary, Conclusion and Outlook

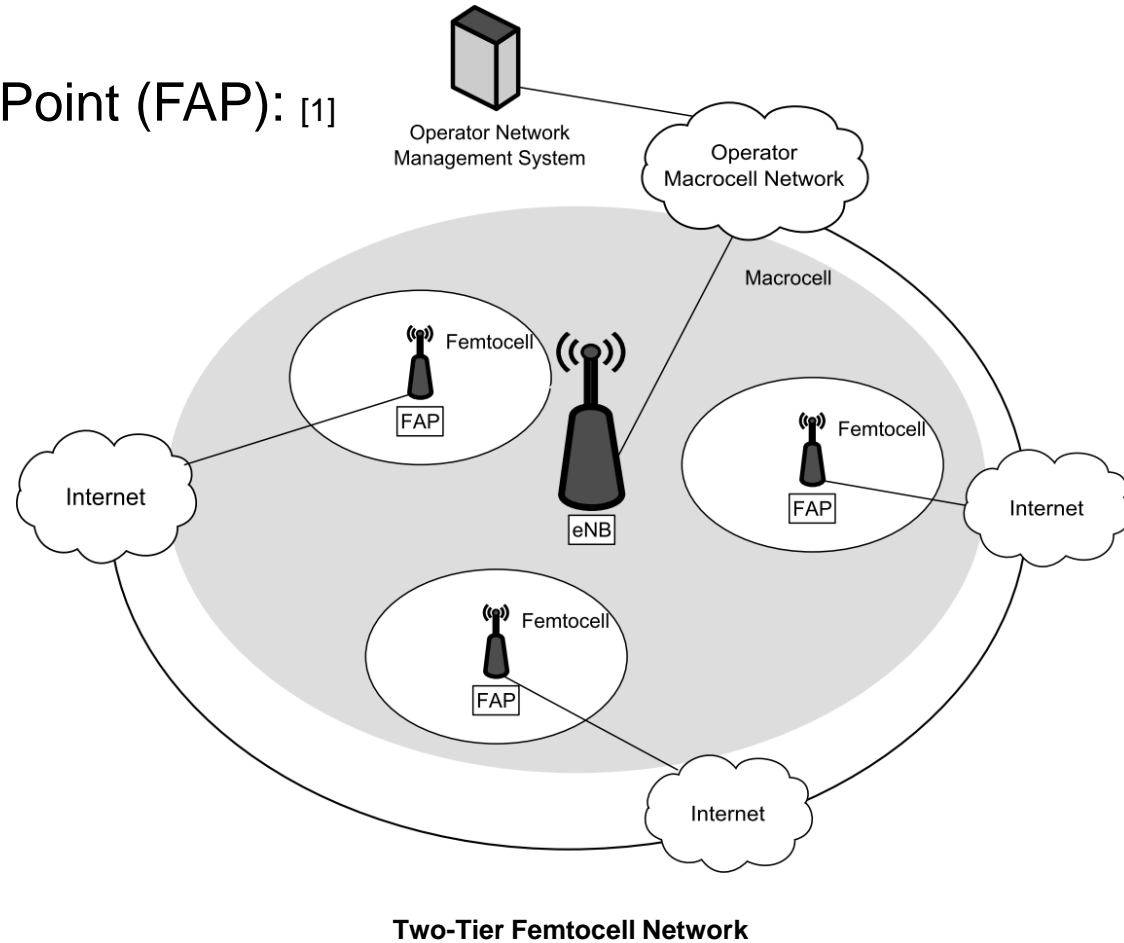
Introduction

Properties of a Femto Access Point (FAP): [1]

- Low power
- Serving small areas
- Operating in licensed spectrum
- Self-organising, self-managing
- IP-Backhaul (DSL)

Standardized by 3GPP:

- UMTS: Home Node B (HNB)
- LTE: Home eNode B (HeNB)



[1] ITU-R WP 5D/195

Motivation

70% of data traffic originates indoors

50% of voice calls are initiated from inside [2]

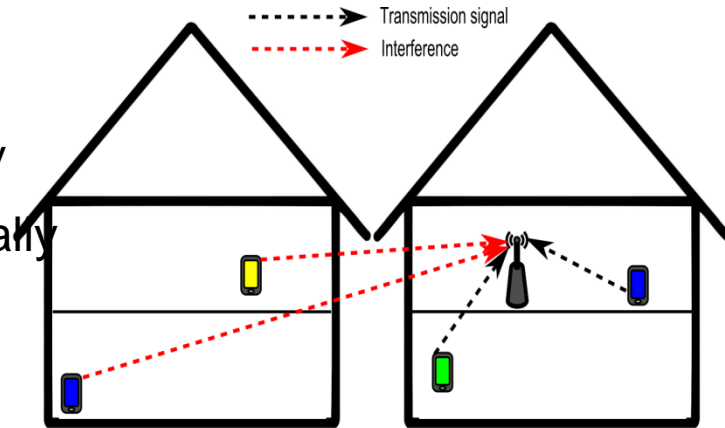
Advantages

- Improved indoor coverage
- Single radio access technology (RAT)
 - Reduced device complexity
 - Seamless handover
- Less power consumption

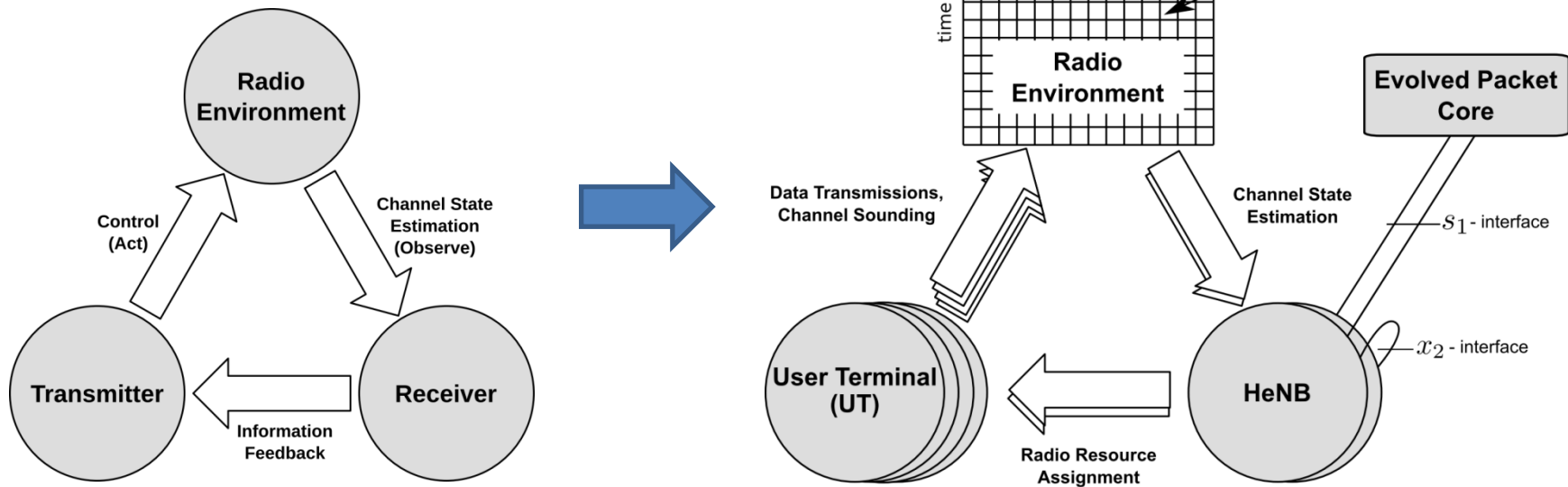
[2] G. Mansfield, "Femtocells in the US Market Business Drivers and Consumer Propositions," Femtocells Europe, ATT, London, U.K., June 2008.

Self-Organizing Networks (SON)

- Number of femtocells can be very large
- Subscriber may switch them on and off frequently
- Operator may not be able to access them physically
- Uplink: Interference depends on UT positioning

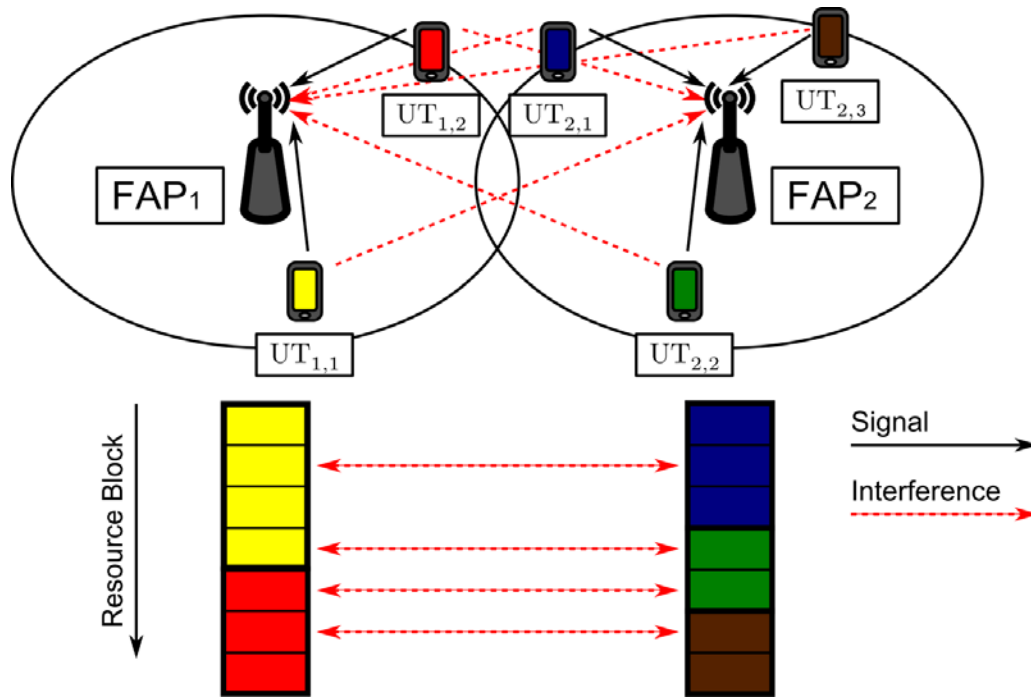


Approach based on Basic Cognitive Cycle [3]



[3] S. Haykin, Cognitive Radio: Brain-Empowered Wireless Communications. IEEE Journal on Selected Areas in Communications, 23(2): 201-220, 2005

Uplink Femtocell Interference Scenario



$UT_{k,j}$ k : Cell index
 j : UT index

Degrees of Freedom:

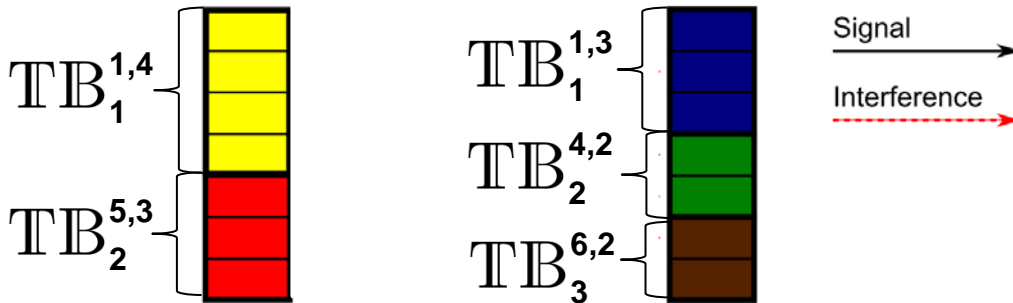
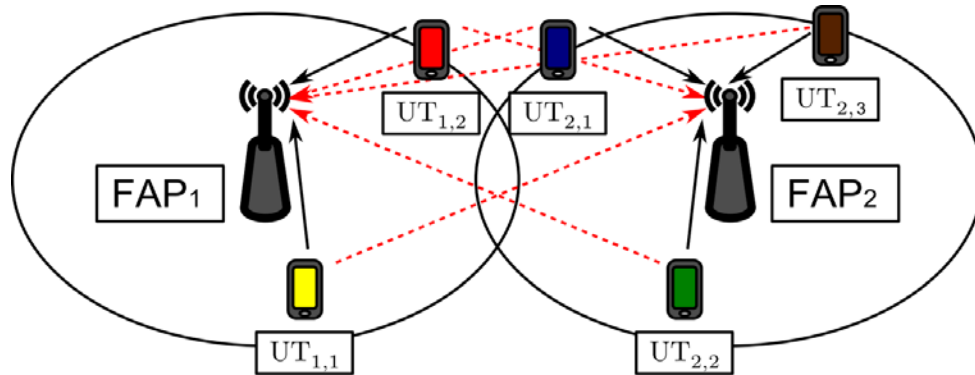
- How many RBs per UT
- Which Transmission Power
- Which Modulation and Coding Scheme (MCS)
- **Which RBs to which UTs**

LTE Constraint:

$$TB_t^{s,l} = \{RB_s, \dots, RB_{s+l}\}$$

s : Start index of first RB
 l : Length of TB

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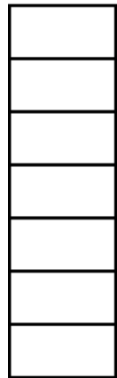
Radio Resource Management Problem

➤ Selecting RBs forming TBs

- Link Adaption
- Assignment of UTs to TBs

Resource Fairness:

Equal amount of RBs in every TB [4]



Number of RBs



Number of UTs

[4] 3GPP. Evolved Universal Terrestrial Radio Access (E-UTRA); Further advancements for E-UTRA physical layer aspects TR 36.814, 2010. V9.0

Radio Resource Management Problem

➤ Selecting RBs forming TBs

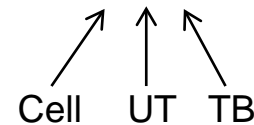
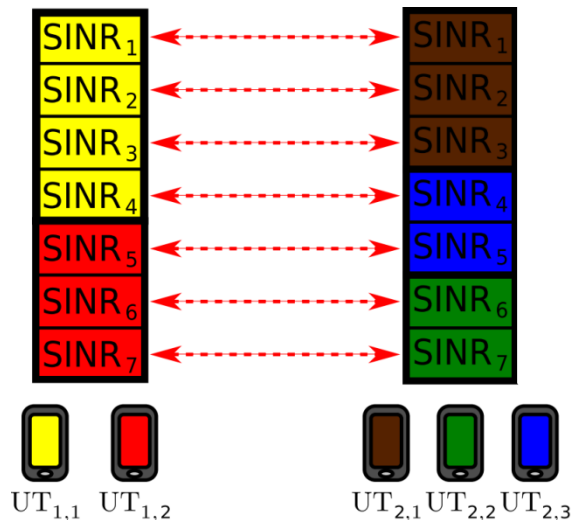
➤ Link Adaption

➤ Assignment of UTs to TBs

Effective SINR [5] for TB with index t:

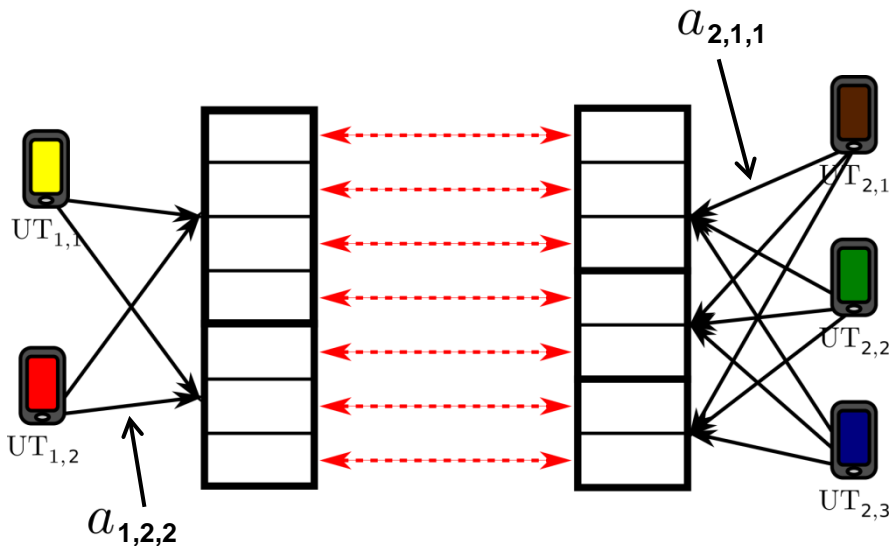
$$\text{SINR}_{eff}(k,j,t) = I^{-1} \left(\frac{1}{|\mathbb{T}\mathbb{B}_t^{s_t,l_t}|} \sum_{n=s_t}^{s_t+l_t-1} I(\text{SINR}_n(k,j)) \right)$$

SINR Mapping with the invertible “information measure” function $I(\cdot)$



Radio Resource Management Problem

- Selecting RBs forming TBs
- Link Adaption
- Assignment of UTs to TBs



Cell Spectral Efficiency:

$$CSE = \frac{1}{KB} \sum_{k=1}^K \sum_{\forall UT_{k,j} \in \mathcal{S}_k} n_{RB} \cdot r(\text{SINR}_{eff}, UT_{k,j})$$

Optimization:

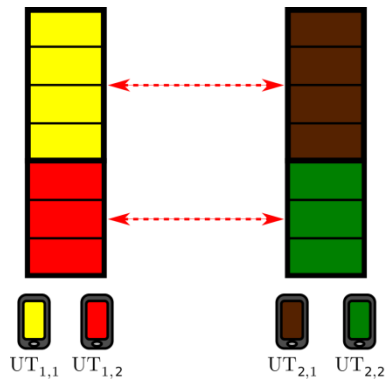
$$\max \sum_{k \in \mathcal{K}} \sum_{j \in \mathcal{S}_k} \sum_{t \in \mathcal{T}_k} a_{k,j,t} \cdot \underbrace{r(\text{SINR}_{eff}(k,j,t))}_{\text{dependent on } a_{k,j,t}}$$

$$a_{k,j,t} = 1 : UT_j \text{ in cell } k \text{ transmits on } TB_t$$

Radio Resource Management Problem

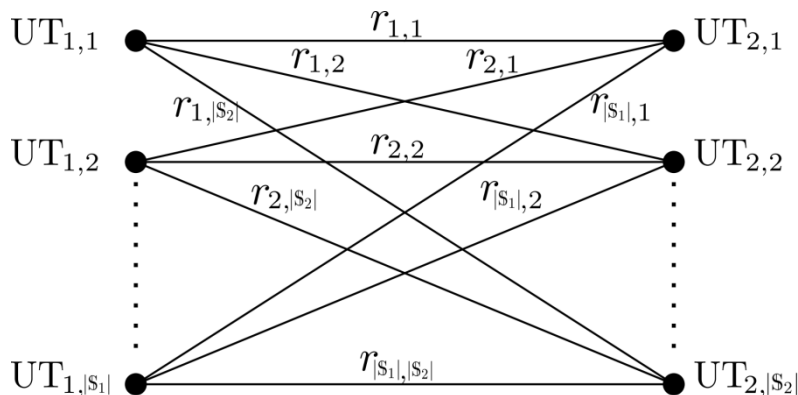
Additional Constraint:

Equal amount of UTs per cell/frame.



Single Interferer from each cell on each $\mathbb{T}B_t$

Reduction to Linear Assignment Problem:



$$\max \sum_{i \in \mathcal{S}_1} \sum_{j \in \mathcal{S}_2} a_{i,j} \cdot r_{i,j}$$

$r_{i,j}$: utility e.g. throughput, SINR

Radio Resource Assignment Algorithms

Two-dimensional algorithms:

UT _{2,1}	$r_{1,1}$	$r_{2,1}$		$r_{ S ,1}$
UT _{2,2}	$r_{1,2}$	$r_{2,2}$		$r_{ S ,2}$
⋮			⋱	
UT _{2, S}	$r_{1, S }$	$r_{2, S }$		$r_{ S , S }$
	UT _{1,1}	UT _{1,2}	⋯	UT _{1, S}

Optimal Algorithm:

Hungarian/Munkres Algorithm
 - Polynomial time $O(n^3)$
 (Operation Research)

Multi-dimensional algorithms:

				UT _{3,5}
				UT _{3,2}
				UT _{3,4}
UT _{2,1}	$r_{1,1,1}$	$r_{2,1,1}$		$r_{ S ,1,1}$
UT _{2,2}	$r_{1,2,1}$	$r_{2,2,1}$		$r_{ S ,2,1}$
⋮			⋱	
UT _{2, S}	$r_{1, S ,1}$	$r_{2, S ,1}$		$r_{ S , S ,1}$
	UT _{1,1}	UT _{1,2}	⋯	UT _{1, S}

No Optimal Algorithm (NP-hard [6]):

Brute Force $O((n!)^{|\mathbb{K}|})$

Number of UTs equals steps

Number of Cells equals dimensions

Heuristics:

Greedy Algorithm $O(n^{|\mathbb{K}|+1})$

Maximum-Regret Algorithm $O(|\mathbb{K}| \cdot n^{|\mathbb{K}|+1})$

[6] M. Garey and D. Johnson, Computers and Intractability, A Guide to NP-completeness, W.H. Freeman and Company: San Francisco, CA, 1979


Greedy Algorithm

Description:

- Locally optimized choices

Example for 2 dimensions:

	UT _{1,1}	UT _{1,2}	UT _{1,3}
UT _{2,1}	7	8	1
UT _{2,2}	9	11	8
UT _{2,3}	7	15	10

-  Step1: Localize maximum uncovered value
- Step2: Cover row and column of maximum value

Solution: {15}9}1}

Sum: 25

Optimal solution: {15,8,7}

Sum: 30

Maximum-Regret

Description:

- Maximizes the regret [7]

Row regret :

$$\Delta_i^r = c_{d_i,i}^r - c_{e_i,i}^r$$

$$c_{d_i,i}^r \geq c_{e_i,i}^r$$

Column regret:

$$\Delta_i^c = c_{i,a_i}^c - c_{i,b_i}^c$$

$$c_{i,a_i}^c \geq c_{i,b_i}^c$$

Example for 2 dimensions:

	UT _{1,1}	UT _{1,2}	UT _{1,3}	Δ^r
UT _{2,1}	7	8	1	
UT _{2,2}	9	11	8	
UT _{2,3}	7	15	10	
Δ^c			7	

Step1: Compute row and column regret Δ^r Δ^c

Step2: Localize maximum uncovered regret value and take the higher one

Step3: Cover row and column of maximum value

Solution: {15,8,7}

Sum: 30

Optimal solution: {15,8,7}

Sum: 30

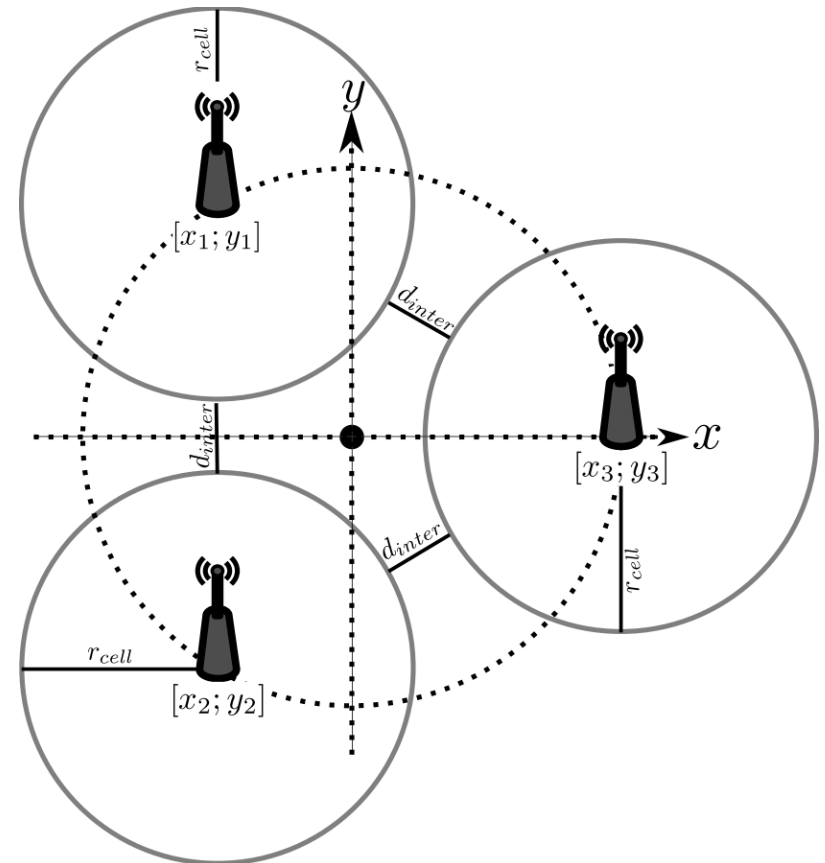
[7] Loomes, G. and Sugden, R. (1982), 'Regret theory: An alternative theory of rational choice under uncertainty', Economic Journal, 92(4), 805–24.

Simulation Scenario

Access to Femtocells: [8]

- **Closed Subscriber Group (CSG) cells:** Only accessible by members of the CSG
- **Hybrid cells:** Accessible as a CSG cell by members of the CSG and as a normal cell by all other UTs.

Parameter	Value
Number of active UTs per cell	2 - 8
Number of cells	2, 3, 4
Inter-cell Edge distance	-100m, -50m, 0m, 50m, 100m



[8] 3GPP TS 32.571 Home Node B (HNB) and Home eNode B (HeNB) management; Type 2 interface concepts and requirements; Release 10

Simulation Parameter

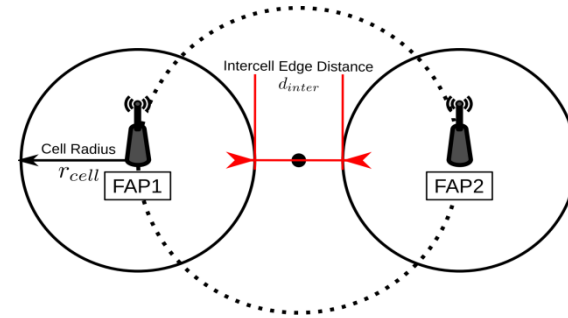
Parameter	Value	Comment
Duplex Scheme	FDD	---
Frequency Band	2500 MHz	---
System Bandwidth	20 MHz (100 RBs)	---
Frame Length	1 ms	---
Pathloss Model	InH NLoS	Guidelines for evaluation [9]
Traffic Model	Full Buffer	---
Adaptive MCS	13 MCSs	---
Transmit Power P_{Tx} UT	4.0 dBm (constant per subchannel)	---
Evaluated Sim. Runtime	1s	---
Cell Radius	100m	---

[9] ITU-R. Guidelines for evaluation of radio interface technologies for IMT-advanced m.2135, 2009.

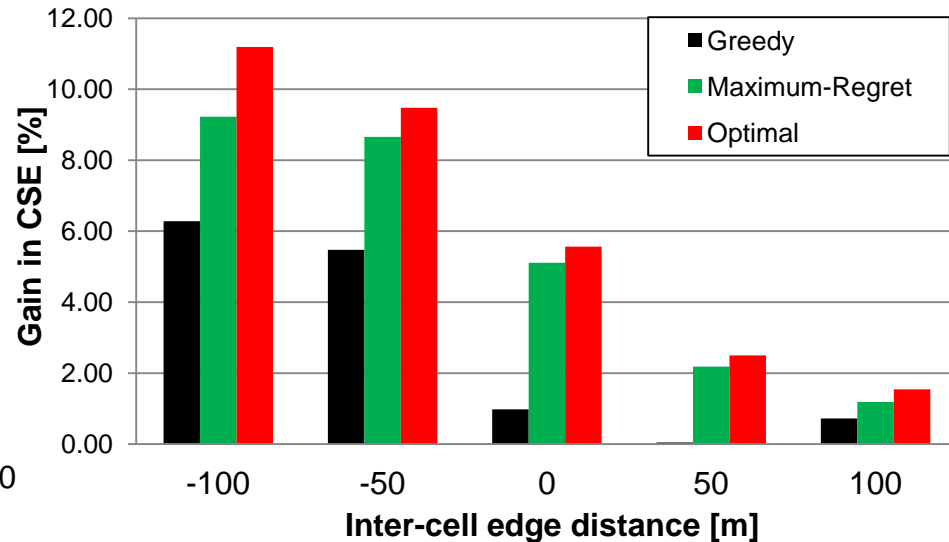
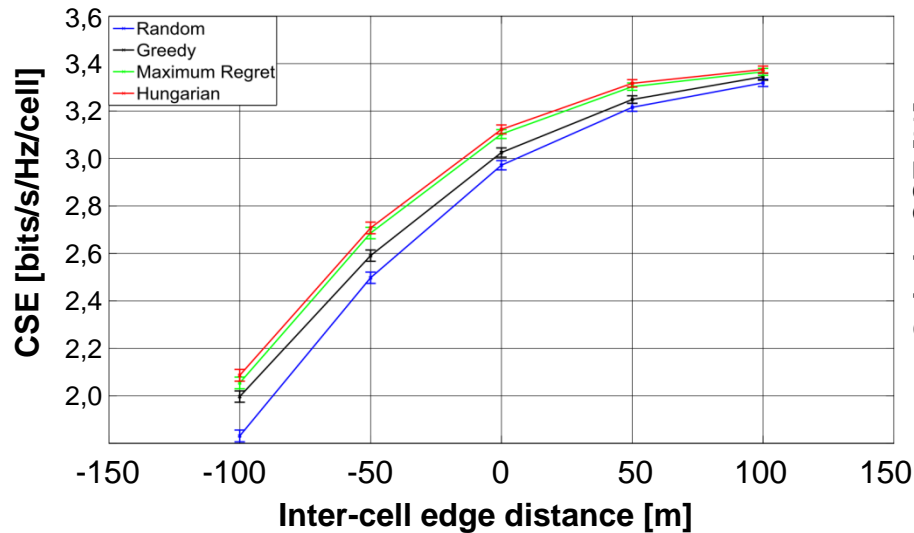
Results: Two Cell Scenario

Comparison of the Strategies to:

- Random (Round Robin in Frequency Domain)



- 2 cells
- 5 UTs per cell



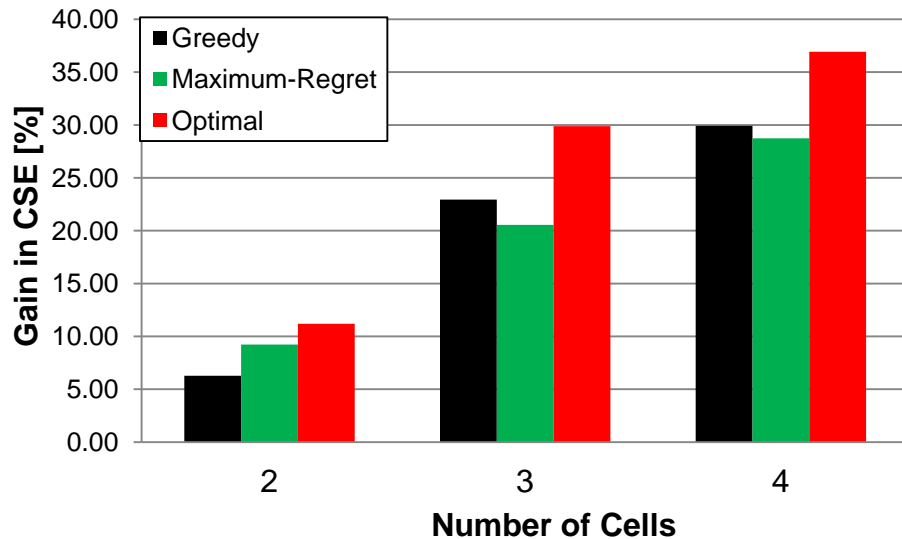
- CSE increases with inter-cell distance
- Strategies merge for increasing inter-cell distance

- Maximum-Regret produces close to optimal results

Results: Multicell

Comparison of the Strategies to:

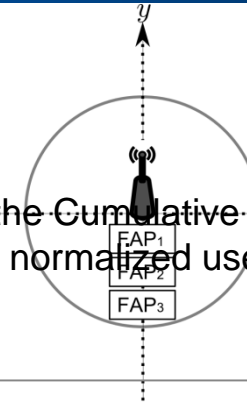
- Random (Round Robin in Frequency Domain)



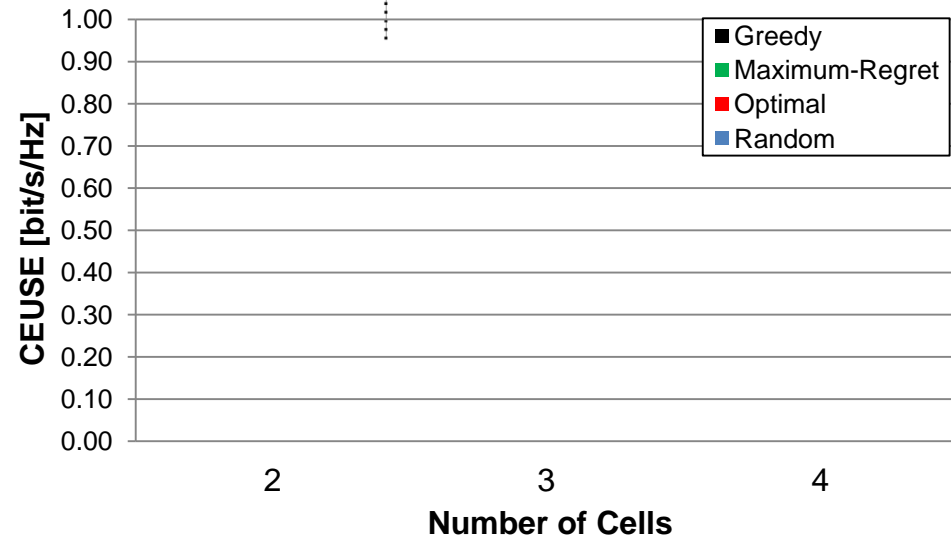
- CSE gain increases with the number of Cells

Fairness (CEUSE)

5% point of the Cumulative Distribution Function (CDF) of the normalized user throughput.



- Spectral Efficiency
- 5 UTs per cell
- $d_{inter} = 100m$

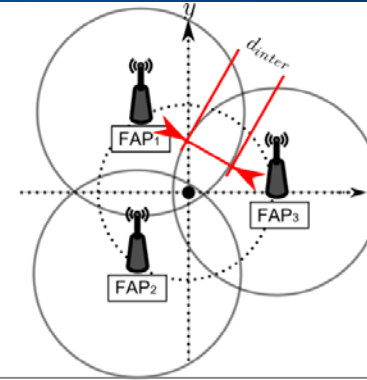


- the worst 5% of users are not served

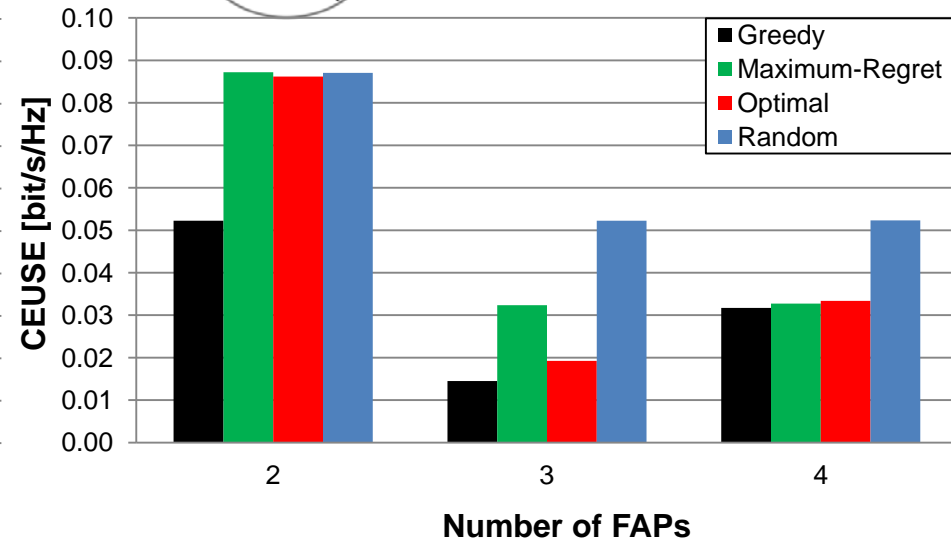
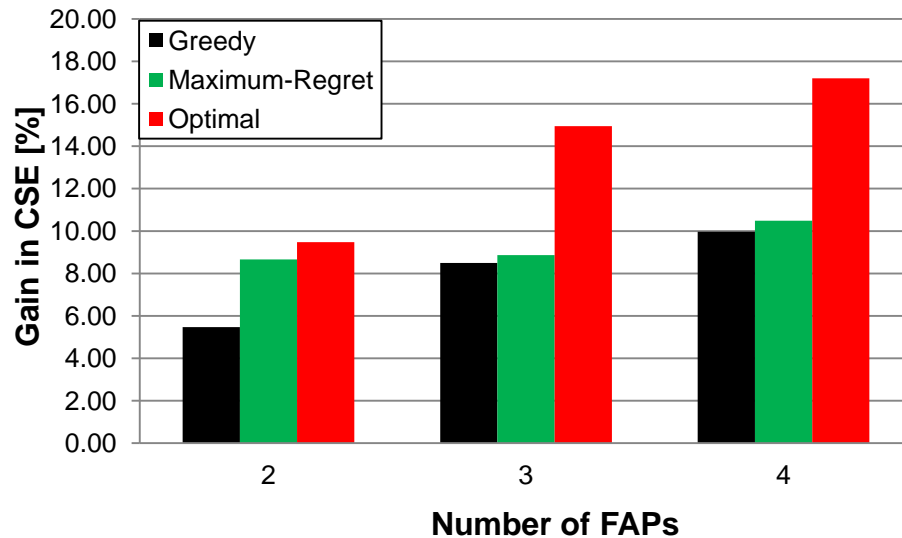
Results

Comparison of the Strategies to:

- Random (Round Robin in Frequency Domain)



- 2, 3, 4 cells
- 5 UTs per cell
- $d_{inter} = -50m$



- CSE gain increases with the number of FAPs

- Cell Edge User suffer from CSE optimization
- Maximum-Regret improves fairness for two and three cells.

Summary and Conclusion

Summary:

- Described the challenges of uplink radio resource management
- Researched and implemented suitable RRM algorithms
- Implemented a central scheduling unit in openWNS Simulator
- Evaluated algorithms in openWNS Simulator

Conclusion:

- Centralized RRM strategies can significantly increase the CSE (up to 37%).
- Achieved CSE gain depends on inter-cell interference power.
- CEUSE suffers from maximizing CSE.

Outlook

Outlook:

- Additional RRM algorithms such as Decomposable Costs [10]
- Same number of UTs per frame.
- Compare against decentralized resource assignment strategies.
- Reduction of Interferer: The assignment of the radio resources can be computed partially with the focus on the strongest interferer.

[10] Approximation algorithms for multi-dimensional assignment problems with decomposable costs. H. Bandelt, Y. Cramab, F. C.R. Spieksma.

Thank you for your attention!

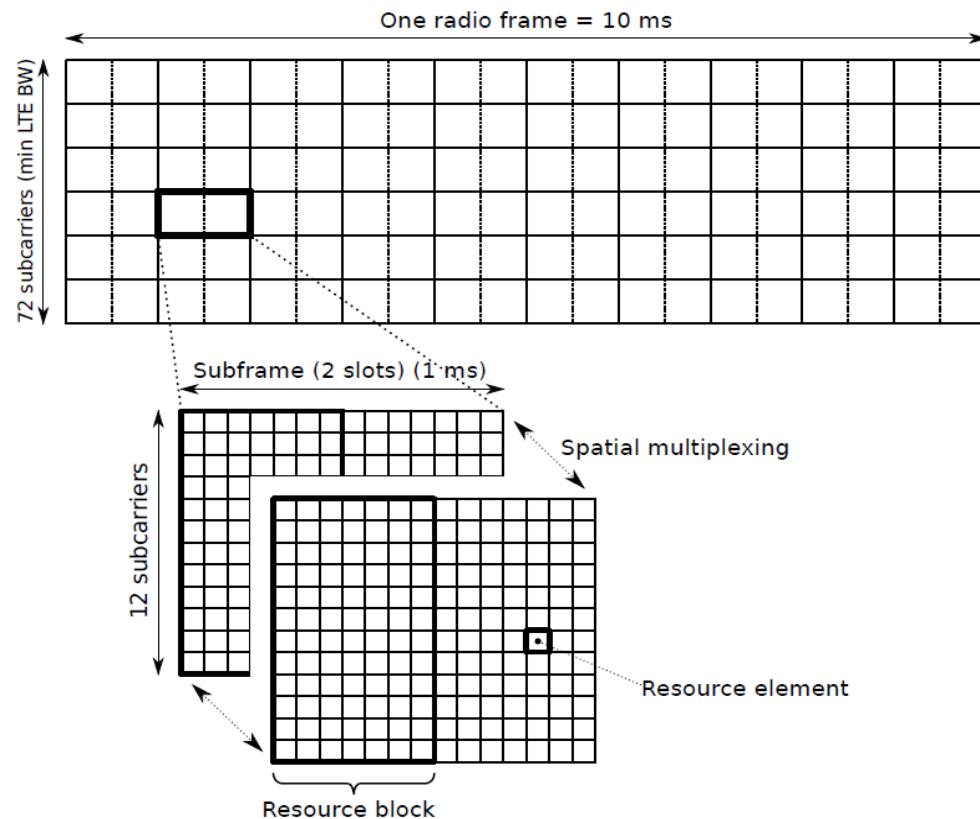
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Fragen und Diskussion

Radio Environment / Resource Structure

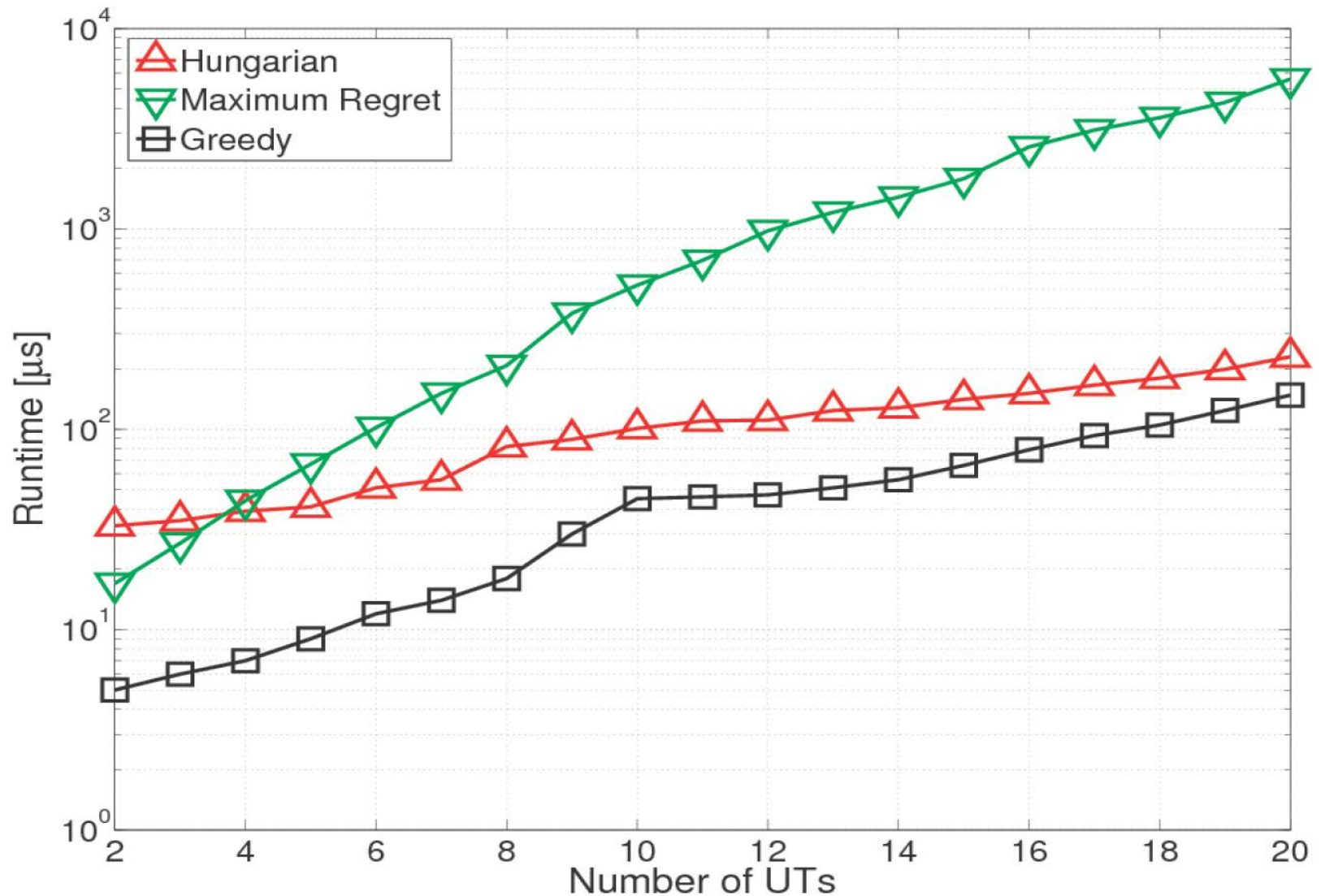
Overview – Introduction – **Motivation** – Problem description – Solution concepts –Workplan



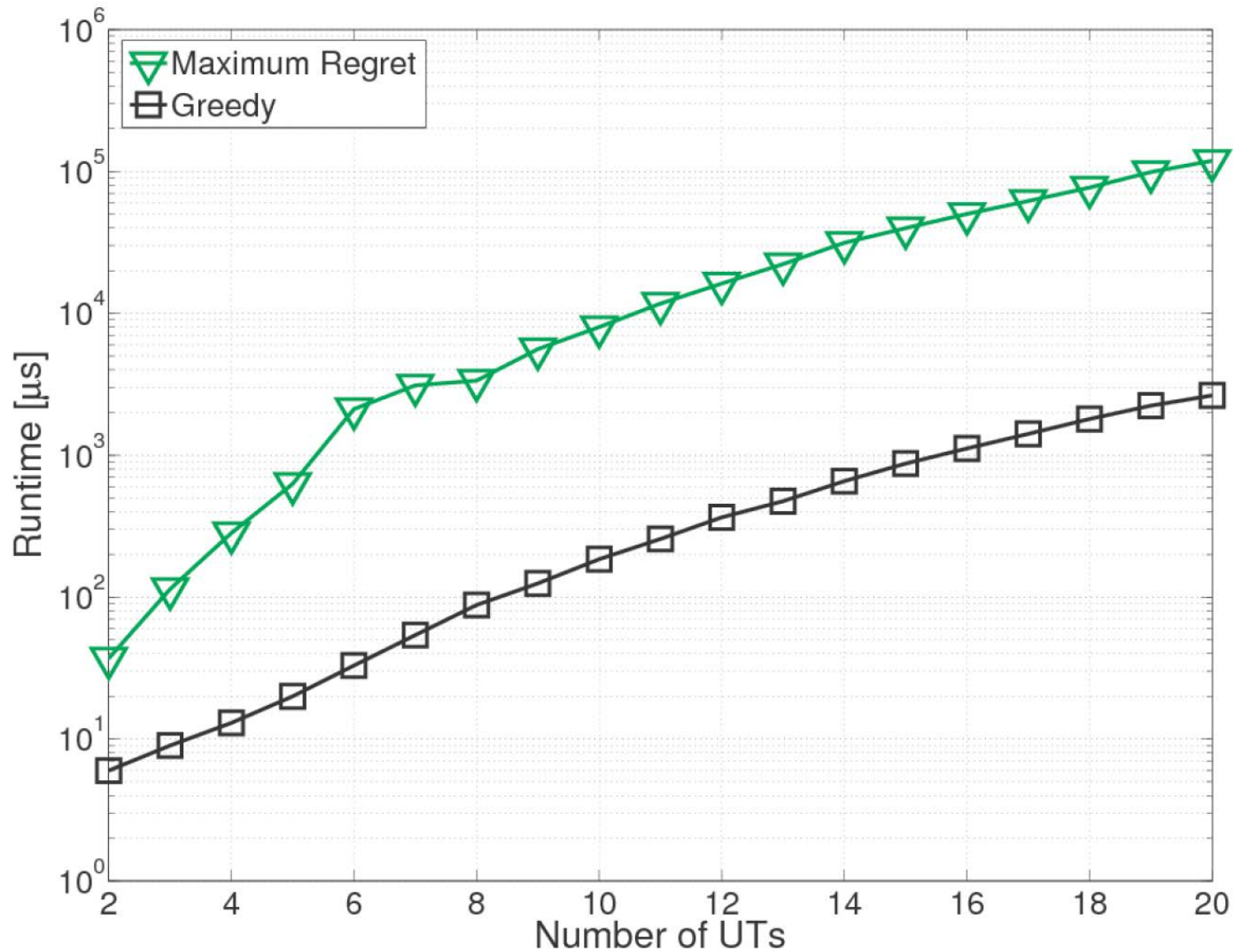
Resource Block (RB)

- In the following the smallest resource of a system is a resource block (RB)
- RBs for each UT are allocated continuously in transport blocks (TBs) (*LTE requirement*)

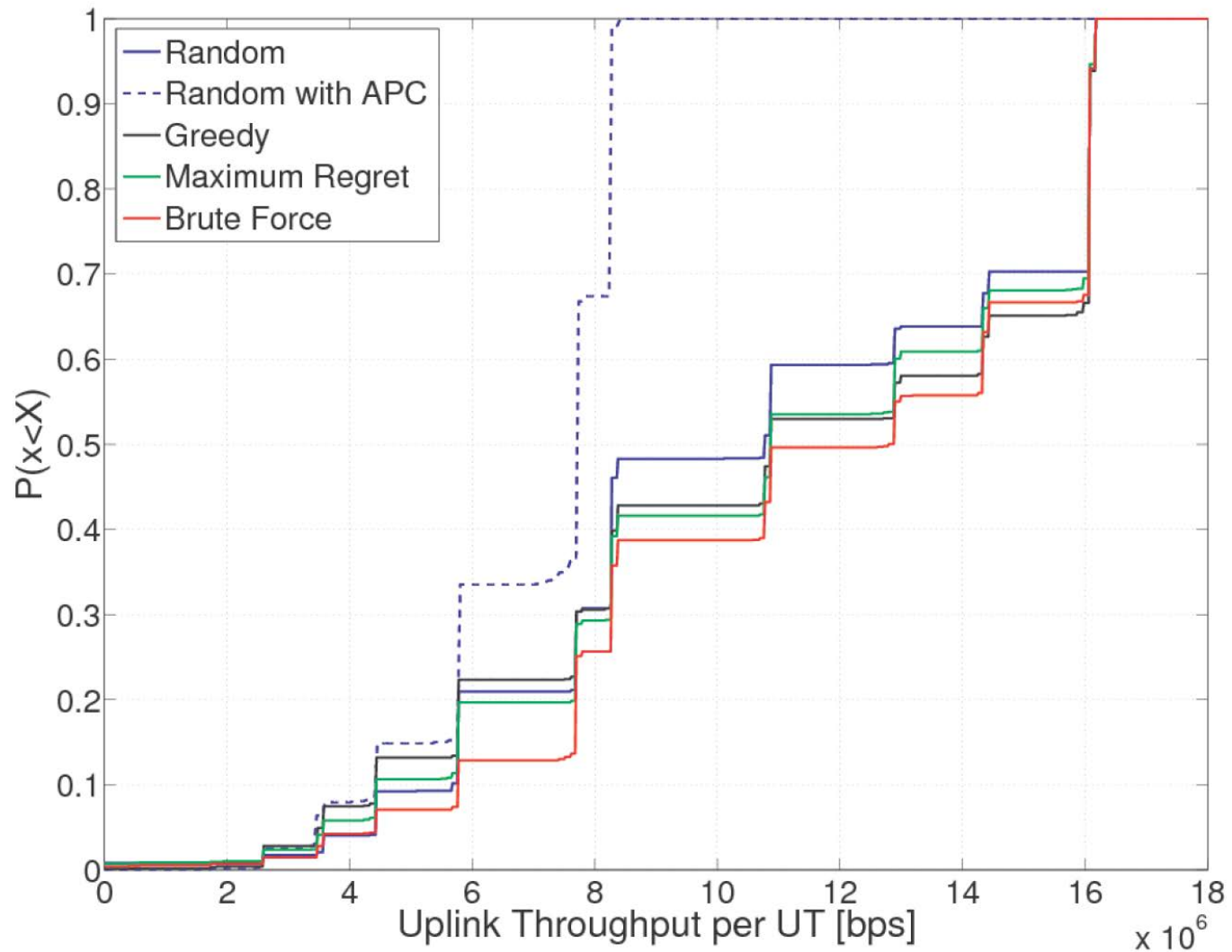
Runtime: 2FAPs



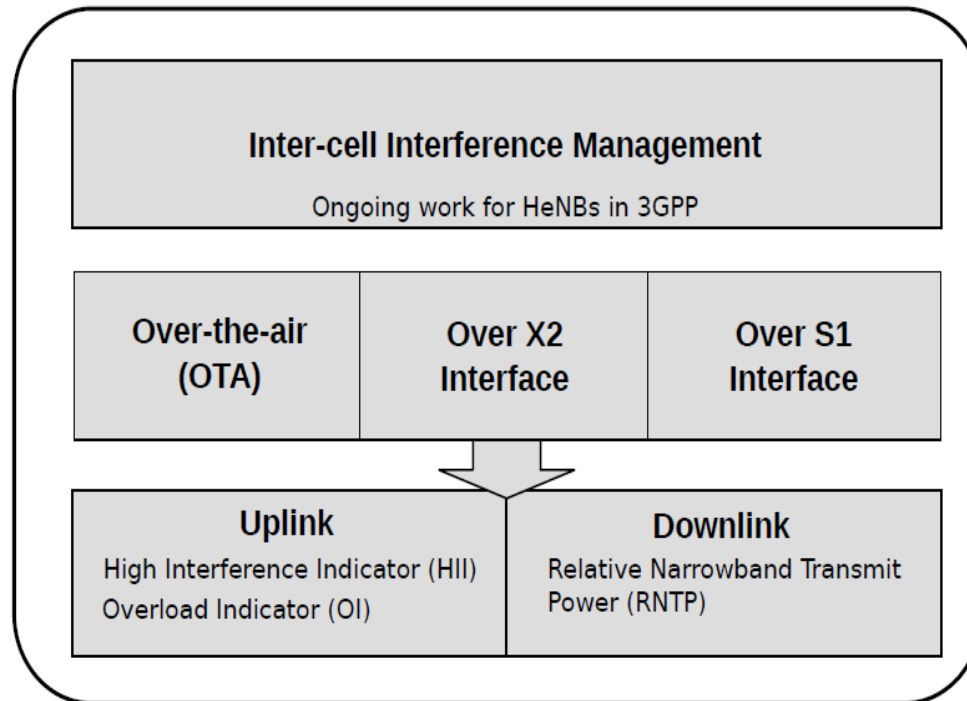
Runtime: 3FAPs



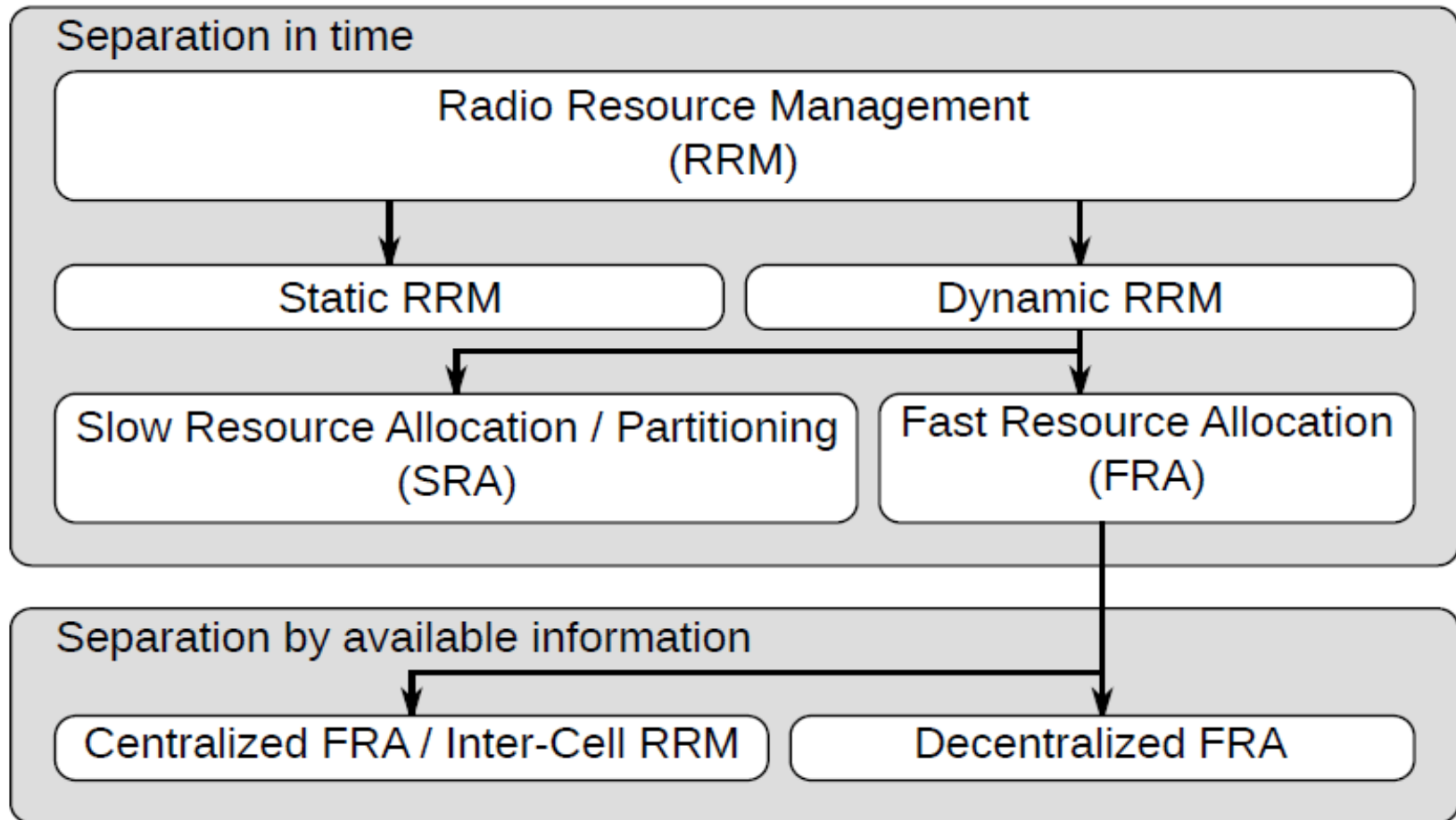
CDF: 3FAPs and 5UTs



Channel State Information



Radio Resource Management

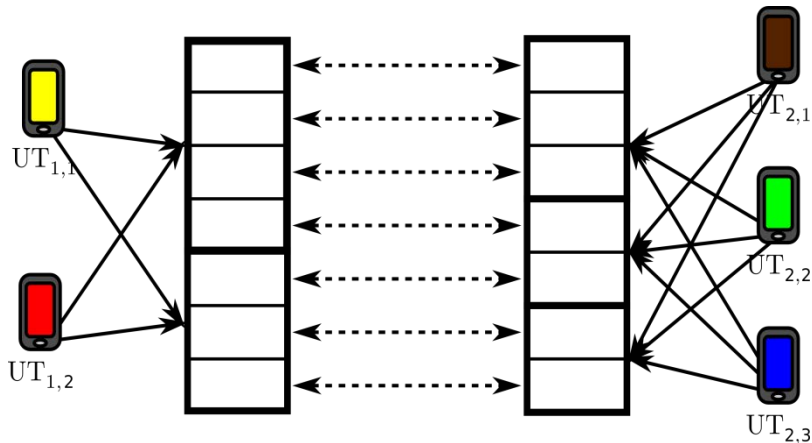


Radio Resource Management Problem

➤ Selecting RBs forming TBs

➤ Link Adaption

➤ Assignment of UTs to TBs



Cell Spectral Efficiency:

$$CSE = \frac{1}{KB} \sum_{k=1}^K \overbrace{\sum_{\forall UT_{k,j} \in \mathbb{S}_k} n_{RB}(UT_{k,j}) \cdot r(\text{SINR}_{eff, UT_{k,j}})}^{\text{Cell Throughput}}$$

Optimization:

$$\max \sum_{k \in \mathbb{K}} \sum_{j \in \mathbb{S}_k} \sum_{t \in \mathbb{T}_k} a_{k,j,t} \cdot r(\text{SINR}_{eff}(k,j,t))$$

Constraints:

$$\sum_{t \in \mathbb{T}_k} a_{k,j,t} = 1 \quad \text{for } j,k \mid j \in \mathbb{S}_k \wedge k \in \mathbb{K}$$

$$\sum_{k \in \mathbb{K}} a_{k,j,t} = 1 \quad \text{for } j,t \mid j \in \mathbb{S}_k \wedge t \in \mathbb{T}_k$$

$$\sum_{j \in \mathbb{S}_k} a_{k,j,t} = 1 \quad \text{for } t,k \mid t \in \mathbb{T}_k \wedge k \in \mathbb{K}$$

$$a_{k,j,t} \in \{0,1\}$$