# SliceTime

### A platform for accurate and scalable network emulation



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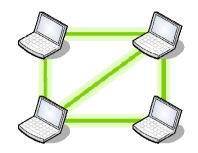
Communication and Distributed Systems

Aachen, 14.7.2011



### **Motivation**

### How to evaluate networking software at large scale?



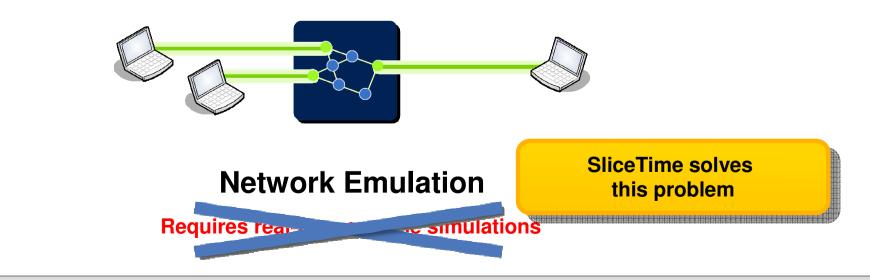


**Network Simulation** 

### **Network Testbeds**

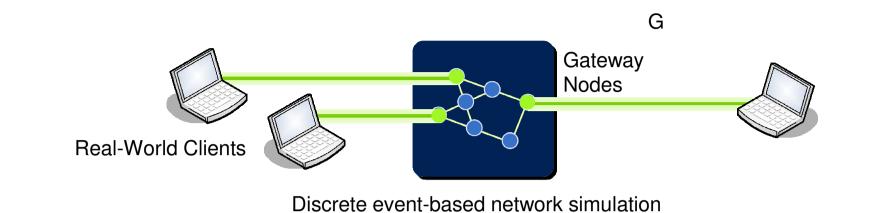
**Drawbacks: Scalability & Cost** 

Models instead of software, no operating system...



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### **Network Emulation**



# Real-World clients

Execute communications software & operating system

### Discrete event-based network simulator

- Models interconnecting network
- Examples: ns-2, ns-3, OMNeT++
- ► Also provides simulated hosts → scalability
- Simulated environment: virtual mobility, radio propagation...

### **Network Emulation: Timing**

# Different timing concepts

- Network simulation: series of discrete events
- Real-world clients: continuous wall-clock time

# Current common solution

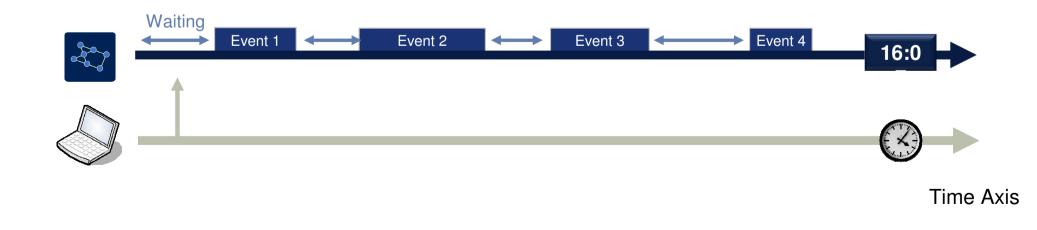
- Pin simulation events to wall-clock time
- Wait between events

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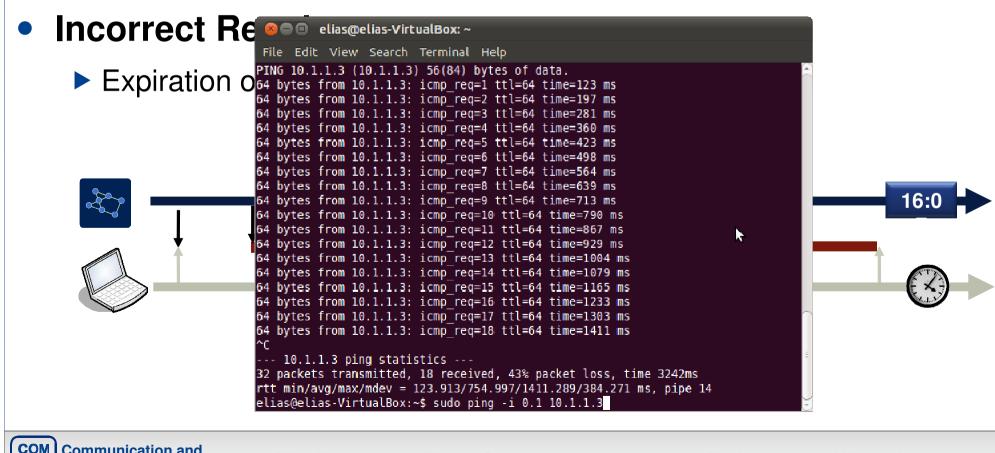


### **Time Drifting Issue**

### • Problem: Many Simulations are not real-time capable

- Computationally complex models
- Many simulated nodes

# Simulation is overloaded → time drift

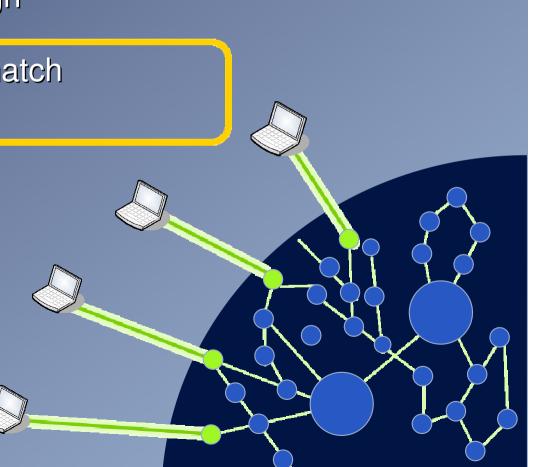


How can time drifting be prevented to enable largescale and complex network emulation scenarios?

Two options:

1. Make the simulation fast enough

2.Slow down the real clients to match the simulation's speed



### **Requirements for SliceTime**

# **1.** Tight synchronization of clients and simulation

Limit drifting to 1ms or less (for WAN scenarios)

# 2. Need to slow down real-world software clients

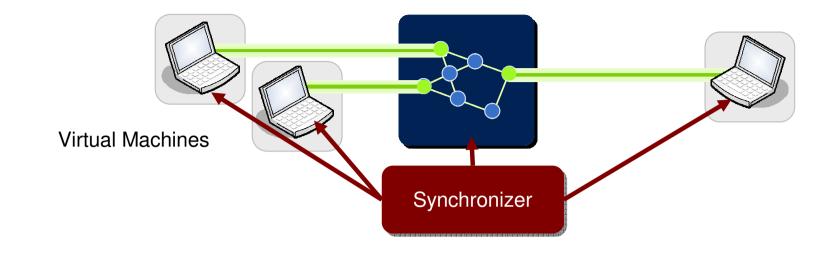
- Unmodified communications software
- Legacy operating systems (Linux or Windows)
- Slow down must be transparent to the clients

   provision of virtual time

# 3. Little overhead due to synchronization

- Additional run-time
- Additional delays or measurement artifacts

# **SliceTime: A Synchronized Network Emulation platform**



# Synchronizer

- Synchronization algorithm aligns execution of clients and simulation
- Virtual machines provide needed level of control
  - Control over run-time behavior
  - ► Full control over system context/timers → provision of virtual continuous time

### **Synchronization Algorithm**

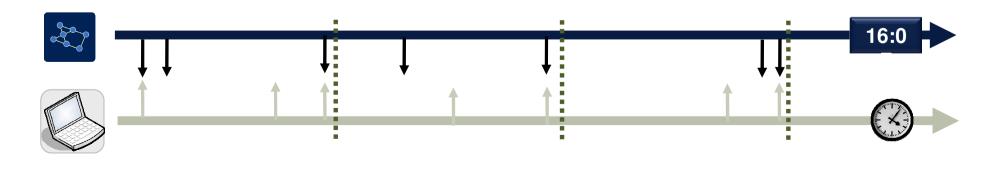
# Goal: Limit time drifting

- No assumptions about future run-time behavior
- No snapshotting & rollbacks

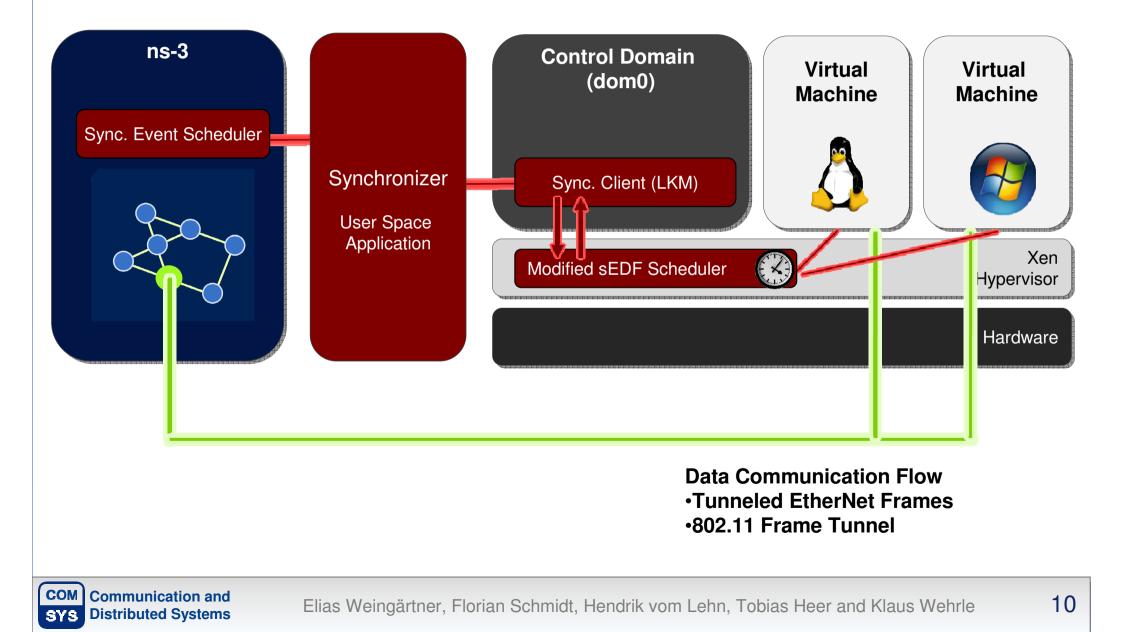
# Barrier Algorithm

- Assign slices of run-time
- Blocking at end of time slice
- Clients notify synchronizer after they have finished

### Synchronization accuracy corresponds to time slice size



### **SliceTime Implementation**

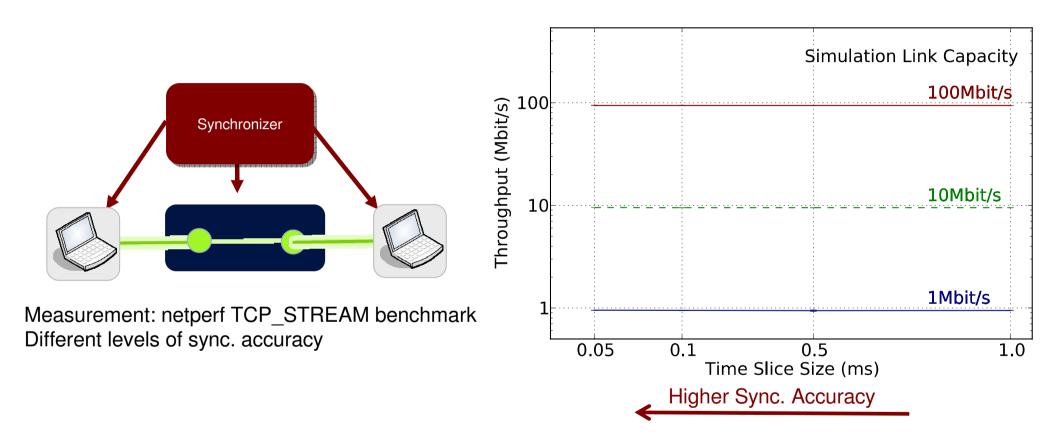


# **Evaluation**

How accurate is SliceTime? How much overhead is caused by the synchronization?

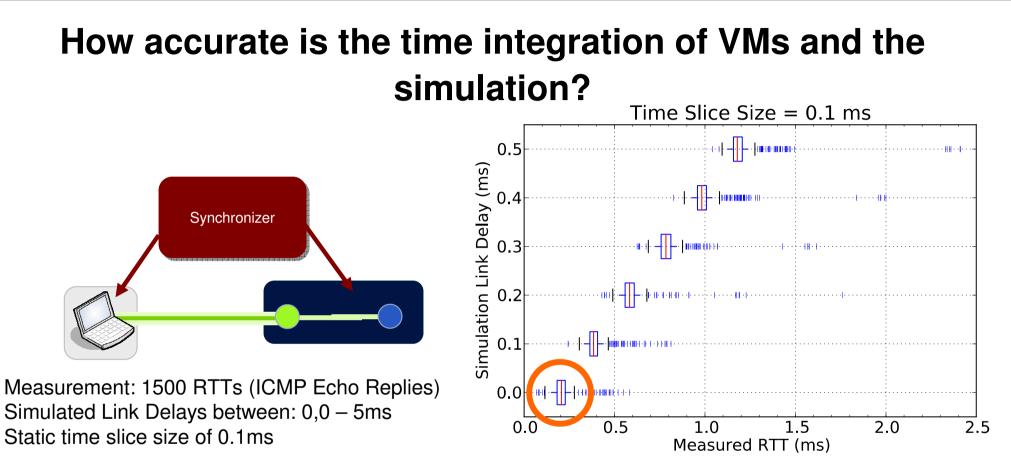
### **Network Throughput**

### How is network throughput affected by time slice size?



Perceived bandwidth is invariant to time slice size

# **Evaluation: SliceTime Timing**



# • If no simulation delay is present $\rightarrow$ RTTs around ~ 0.2ms

Base delay: Time needed for data exchange between VM & sync

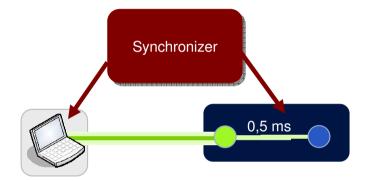
RTT distributions shifted by twice simulation delay

CON

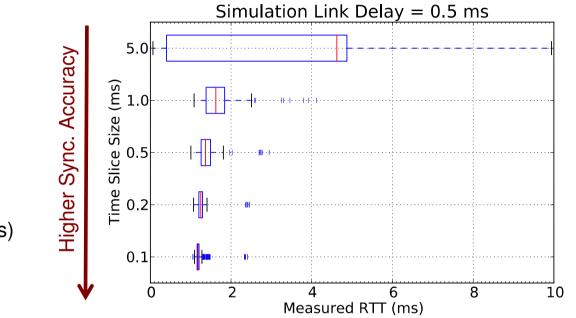
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# **Evaluation: SliceTime Timing**

### How do different time slice sizes influence the results?



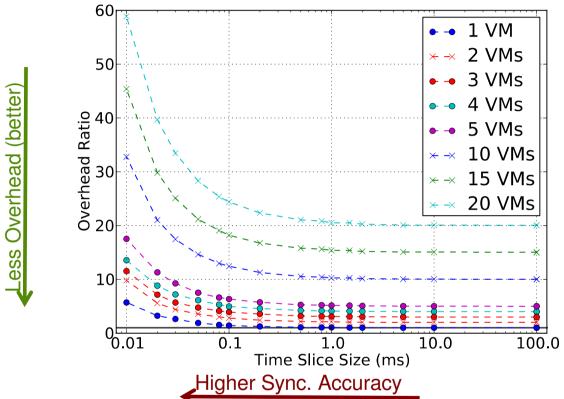
Measurement: 1500 RTTs (ICMP Echo Replies) Variation: Time Slice Sizes



• RTT distributions converge to base delay for smaller time slices (higher accuracies)

### **Synchronization Overhead**





Synchronization introduces additional run-time overhead

- Less than 5% for time slices > 0,5ms
- Linear in the number of VMs

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### Conclusion

- SliceTime allows network emulation scenarios with network simulations of any run time behavior
- SliceTime is accurate regarding timing and throughput
- SliceTime is resource efficient
  - Low overhead even for time slices less 1ms
  - Saves physical hardware resources in comparison to real test beds
- SliceTime is open source
  - Get it at <u>http://www.comsys.rwth-aachen.de/projects/slicetime</u>
- SliceTime extends the applicability of network emulation

# **Questions?**

### Implementation

Two main tasks:

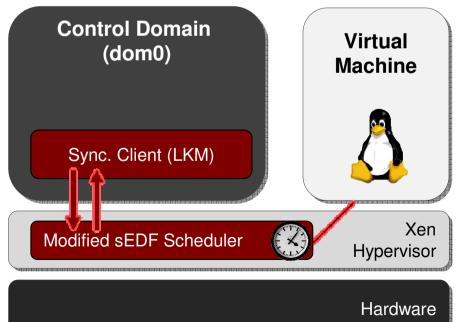
Execute VMs & simulation for exact time slice duration
 Provide VMs with illusion of gapless virtual time

# Synchronization Client

► Linux Kernel Module  $\rightarrow$  save context switches

# Modified sEDF scheduler

- Execute Xen domains for time slice duration
  - Extra scheduling queue for synchronized domains
  - Self-correction mechanism to overcome misattribution of run-time
- Virtualizes time progression for synchronized domains
  - Calculates delta values for timers and clock sources



### **Network Simulation**

# Synchronized Event scheduler

- Synchronizes any ns-3 simulation with synchronizer/VMs
- Checks if next event in queue resides in current time slice

### Different ns-3 extensions

- $\blacktriangleright$  Tunnel protocol  $\rightarrow$  data exchange with VMs
- WiFi emulation extensions
  - Provides VMs with wireless networking interface
  - Interface is intergrated with 802.11 model of ns-3



### Synchronizer

### • Implements barrier synchronization algorithm

- Assignment of time slices
- Synchronizes multiple VMs with multiple simulations

# User-space application

- Can run on VM, simulation slave or dedicated host
- Lightweight signaling protocol

# • VMs and simulations may join sync. dynamically

Allows VM bootstrapping out of synchronization



### Can SliceTime ease the evaluation of networking software?

### AODV Experiment (Gray et al, 2003)

- •33 laptops running AODV
  •40 people carrying them around (on an athletic field)
  •Random UDP traffic
- •Laptops log traffic + position (GPS)
  - Logs available at CRAWDAD

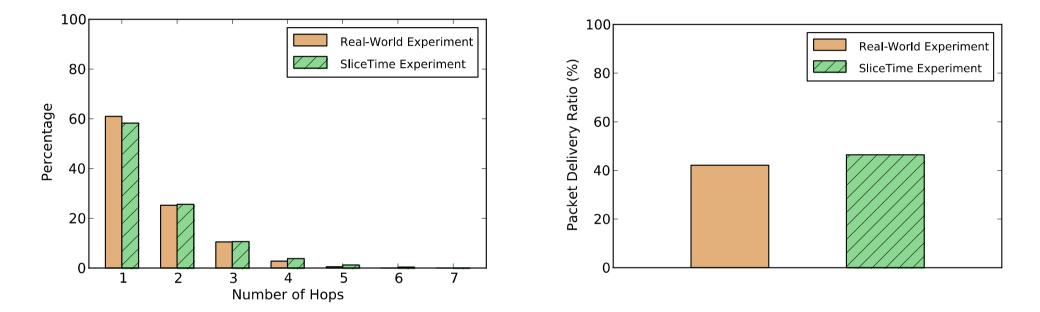
### The SliceTime equivalent

•33 Xen HVM domains / AODV
•SliceTime 802.11 extensions
•1 physical PC
•Ns-3 mobility model based on GPS traces
•Traffic generator



### **Reproducing the AODV experiment by Gray**

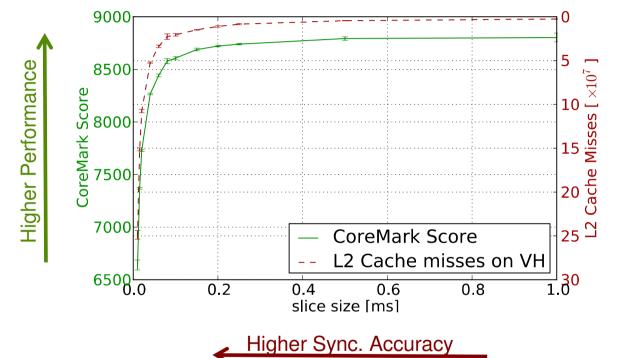
### How do the results compare?



- SliceTime produces results close to real-world measurements
- Always differences due to real-world/simulation disparity

### **CPU Performance Impact**

# How about the CPU performance? Doesn't the synchronization cause artifacts?



CoreMark score decreases for small time slices

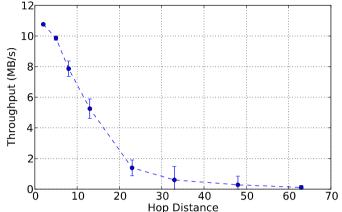
- Almost no impact for slices greater than 0.1ms
- Explanation: More L2 cache misses

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### **SliceTime Simulation scalability**





### Setup: 15000 simulated nodes (60 stars with 250 nodes)

- Exchange data blocks among each other using HTTP
- Executes~15 times slower than real-time
- 1 VM attached to backbone

### HTTP perormance measured with curl

Expected result

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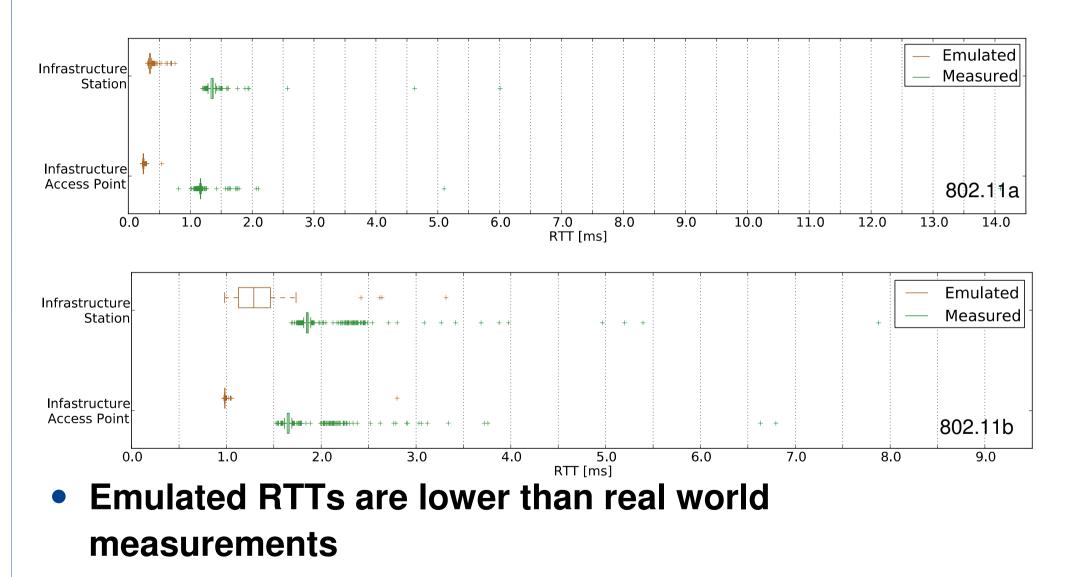
### 802.11 Round Trip Times

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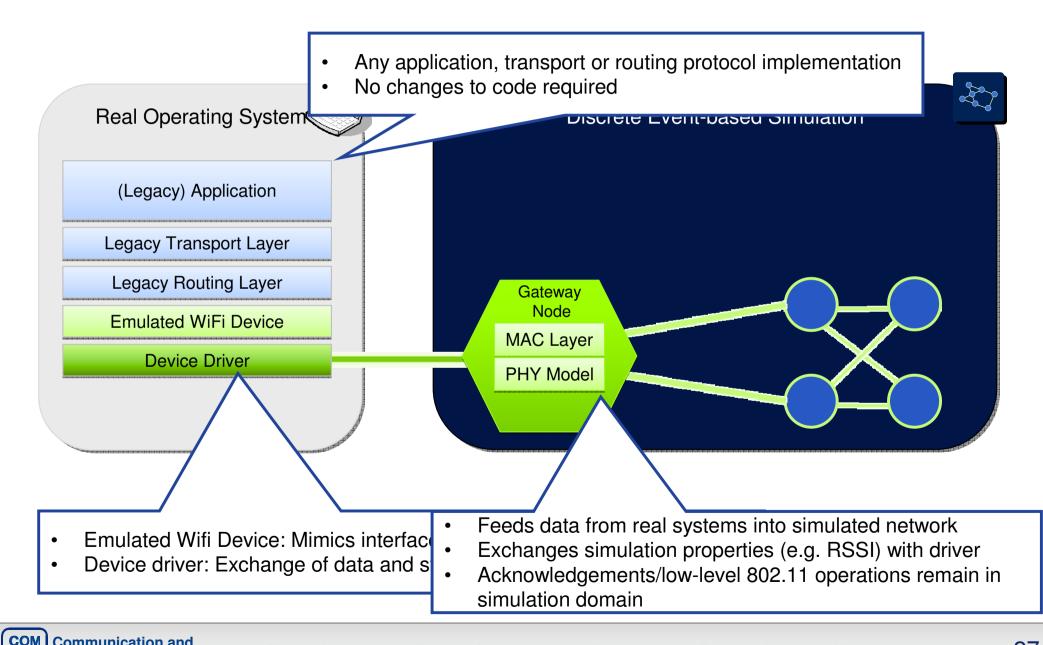




ns-3 only approximations for link-level delays; no system delays

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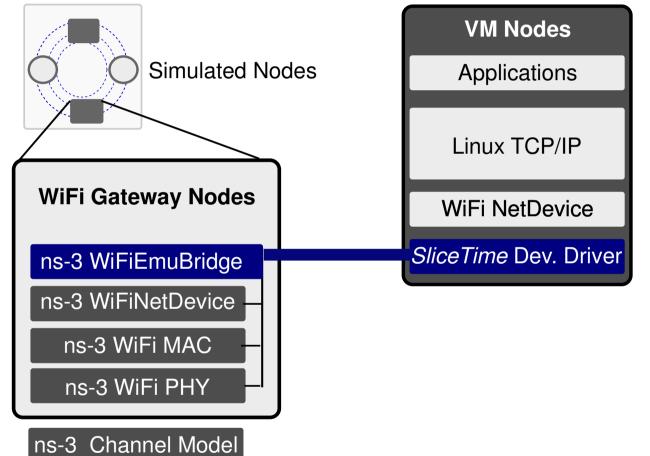
### **Device Driver-enabled Wireless Network Emulation**



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### **SliceTime WiFi extensions**







# **Legacy Applications**

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### • Wireshark for live monitoring of simulated WiFi networks

Inspection of low-level 802.11 properties using Radiotap headers

# **Legacy Applications**

"Kismet Sort View Mindows		Kismet_2
Network-1 A N 3 135 4K		<u>N13060_2</u>
Network-3 A N 3 214 4K		Elapsed
BSSID: 00:00:00:00:00:06 Last seen: Nov 6 19:25:27 Crypt: None Manuf: Unknown		00:08.40
Network-4 A N 7 526 5K		vv+v0++v
Network-5 A N 7 483 4K		Networks
Network-6 AN 7 450 4K		9
Network-7 AN 9 187 4K		Č.
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# • Kismet being executed in simulated network

Allows the execution of unmodified legacy applications that make use of Linux Wireless Extensions