

# A Game Theoretic Approach to Load Balancing in Cellular Radio Networks

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# OUTLINE

## MOTIVATION

## BACKGROUND

- General Definitions

- Load Balancing Algorithm

## METHODOLOGY

- Game and Players Definitions

- Utility Function

- Strategy Played By The Underloaded Cell

- Strategy Played By The Overloaded Cell

## RESULTS

# MOTIVATION

- ▶ The amount of load that the overloaded/underloaded cell should offload/accept might be vendor-specific
- ▶ The algorithm calculating the values of these load runs in the eNodeB
- ▶ The aim is to study the effect on the overall network performance if each cell seeks to maximize its own benefit in a non-cooperative manner

# GENERAL DEFINITIONS

- ▶ Each user  $u$  in the mobile network has a load defined by

$$\kappa_u = \frac{D_u}{R(\text{SINR}_u)N_{\text{total}}}$$

where  $D_u$  is the data rate requirement of the user,  $R(\text{SINR}_u)$  is its corresponding throughput per Physical Resource Block (PRB) and  $N_{\text{total}}$  is the total number of PRBs

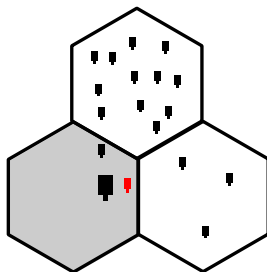
- ▶ The load of a cell  $i$  is denoted by

$$\rho_i = \sum_{u|X(u)=i} \kappa_u \geq 0$$

where  $X(\cdot)$  is the connection function assigning a user  $u$  to a cell

- ▶ If  $\rho_i \leq 1$ , the cell is underloaded
- ▶ If  $\rho_i > 1$ , the cell is overloaded

# LOAD BALANCING ALGORITHM



The overloaded cell having  $\rho > 1$  (Gray) generates the list of candidates to be handed over

- ▶ Users having small link imbalances
- ▶ Users having small data rate
- ▶ Spare capacities in the neighbor underloaded cells

# GAME AND PLAYERS DEFINITIONS

- ▶ **Players:**  
The overloaded and the neighbor underloaded cells
- ▶ **Game:**
  - ▶ The underloaded cells signal to the overloaded cell the amount of load that they are willing to accept
  - ▶ The overloaded cell should offload to each underloaded neighbor cell a certain amount of users having a total load less or equal to the signaled one
- ▶ **Strategies:**  
Each player decides on the amount of load to accept (underloaded cells) or to offload (overloaded cell)

# UTILITY FUNCTION

- ▶ If each player is considered rational, he will decide on a strategy that maximizes his payoff
- ▶ In our game, the utility function can be expressed by the number of satisfied users in the cell
  - ▶ More capacity usage
  - ▶ More income resulting from data rate charging

# UTILITY FUNCTION

- ▶ For the underloaded cell  $i$  having a load  $\rho_i < 1$  and  $U_i$  users

$$\text{utility}_i = \begin{cases} U_i + x_i, & \text{if } 0 \leq \tilde{\rho}_i \leq 1 \\ \left\lfloor \frac{(U_i + x_i)}{\rho_i + \sum_{j=1}^{x_i} \tilde{\kappa}_j} \right\rfloor, & \text{Otherwise} \end{cases}$$

where  $x_i$  is the number of handed over users from the overloaded cell to the underloaded cell  $i$  and

$\tilde{\rho}_i = \rho_i + \sum_{j=1}^{x_i} \tilde{\kappa}_j$  is the load of the cell after load balancing

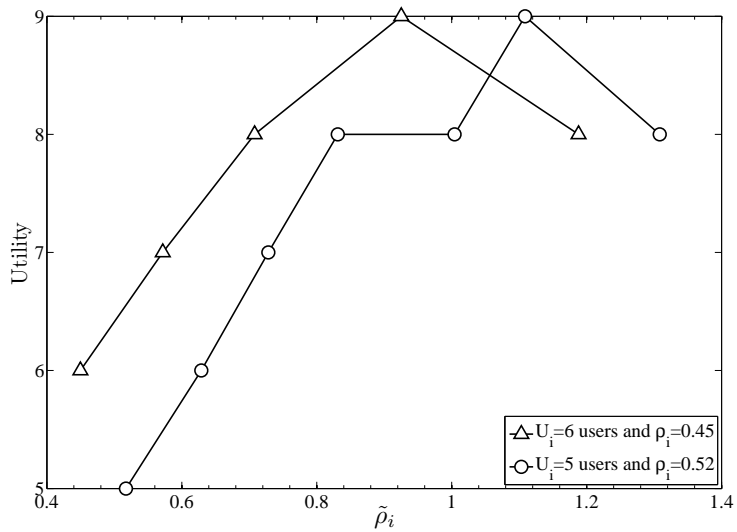
- ▶  $\tilde{\kappa}_j$  is the estimated load of the  $j^{\text{th}}$  handed over user in the target cell  $i$



# UTILITY FUNCTION

$$\text{utility}_i = \begin{cases} U_i + x_i, & \text{if } 0 \leq \tilde{\rho}_i \leq 1 \\ \left\lfloor \frac{U_i + x_i}{\rho_i + \sum_{j=1}^{x_i} \tilde{\kappa}_j} \right\rfloor, & \text{Otherwise} \end{cases}$$

- ▶ The underloaded cell should decide and signal  $y_i$  to the overloaded cell
- ▶ The overloaded cell offload  $\sum_{j=1}^{x_i} \tilde{\kappa}_j \leq y_i$
- ▶ Maximization of the utility function depends on  $x_i$  and  $\tilde{\kappa}_j$  which are not known to the underloaded cell  $i$
- ▶ Utility increases by  $x_i$  if  $y_i \leq 1 - \rho_i$



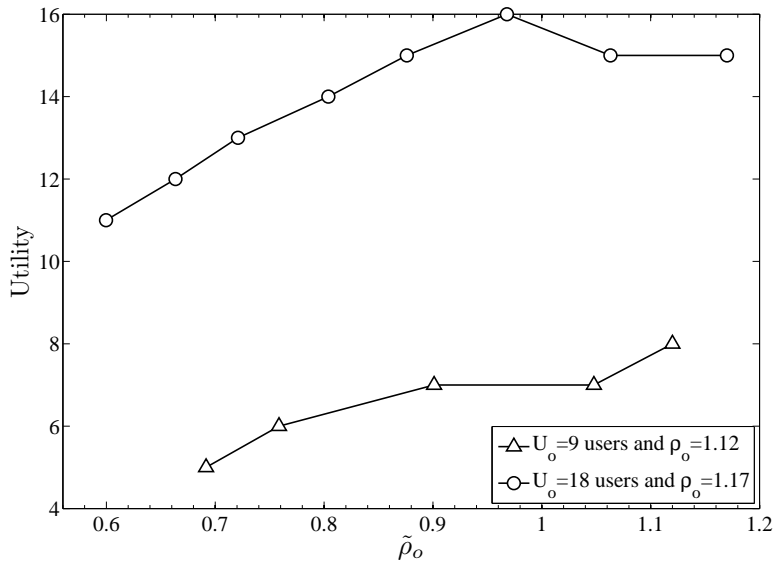
# UTILITY FUNCTION

- ▶ For the overloaded cell having a load  $\rho_o > 1$  and  $U_o$  users

$$\text{utility}_o = \begin{cases} U_o - x_o, & \text{if } 0 \leq \tilde{\rho}_o \leq 1 \\ \left\lfloor \frac{(U_o - x_o)}{\rho_o - \sum_{j=1}^{x_o} \kappa_j} \right\rfloor, & \text{Otherwise} \end{cases}$$

where  $x_o$  is the total number of handed over users from the overloaded cell to the underloaded cells,  $\kappa_j$  is the load of the  $j^{\text{th}}$  handed over user and  $\tilde{\rho}_o = \rho_o - \sum_{j=1}^{x_o} \kappa_j$  is the load of the cell after load balancing

- ▶ The overloaded cell knows the load of each user  $\kappa_j$  and should decide on  $\mathcal{X} = \sum_{j=1}^{x_o} \kappa_j$



# STRATEGY PLAYED BY THE UNDERLOADED CELL

- ▶ The cell signals a load value such that its utility will never decrease and is maximized as much as possible
- ▶ At first, the underloaded cell will accept a load  $y^*$  s.t its utility will not decrease even if 1 user is handed over

$$\left\lfloor \frac{U_i + 1}{\rho_i + y^*} \right\rfloor \geq U_i$$

$$\frac{U_i + 1}{\rho_i + y^*} - U_i \geq 0$$

$$\rho_i + y^* \leq \frac{U_i + 1}{U_i}$$

$$y^* \leq \frac{U_i + 1}{U_i} - \rho_i = y_{\max}^*$$

# STRATEGY PLAYED BY THE UNDERLOADED CELL

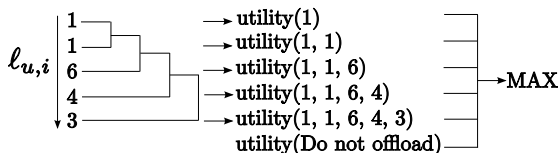
- ▶  $y_{\max}^*$  might be large if the cell is not too much loaded and the underloaded cell might get a small utility gain
- ▶ Hence, the underloaded cell signals  $y_{\max}^*$  only if it is enough loaded, i.e.,  $\rho_i \geq \rho_t = 0.9$

$$y_i = \begin{cases} 1 - \rho_i & \text{If } \rho_i < \rho_t \\ \frac{U_i + 1}{U_i} - \rho_i, & \text{Otherwise} \end{cases}$$

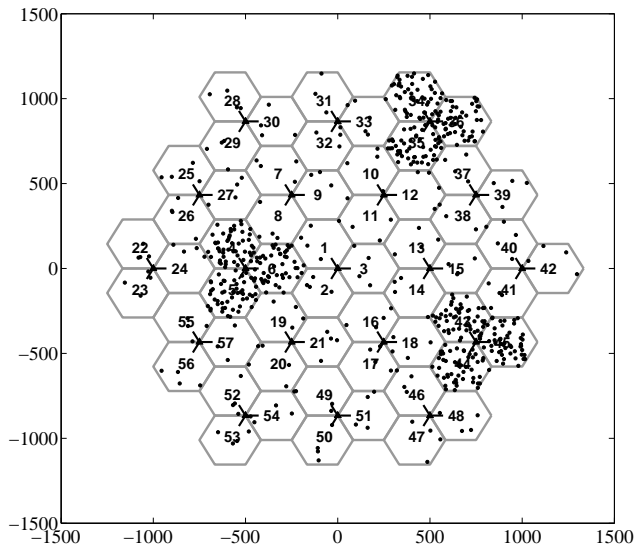
# STRATEGY PLAYED BY THE OVERLOADED CELL

- ▶ It will choose  $\mathcal{X}_{opt}$  that maximizes its utility in a straight forward manner without exceeding the signaled load of each cell

UEs with target cell  $i$

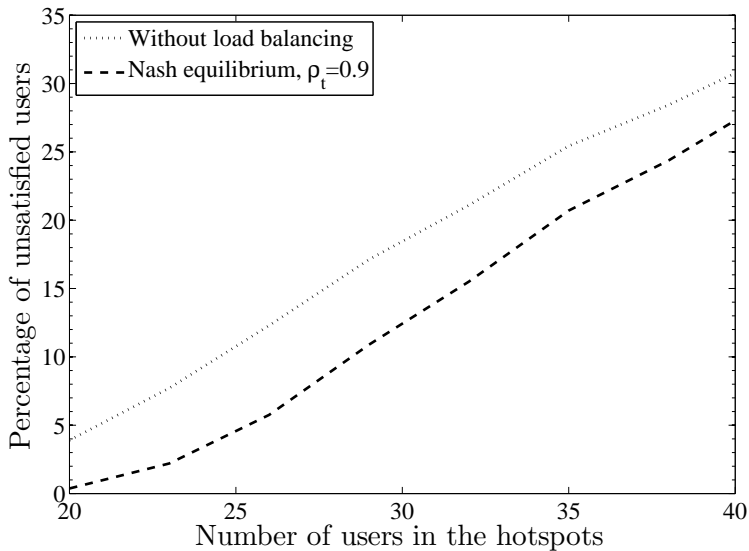


# NETWORK LAYOUT





# SNAPSHOT EVALUATION



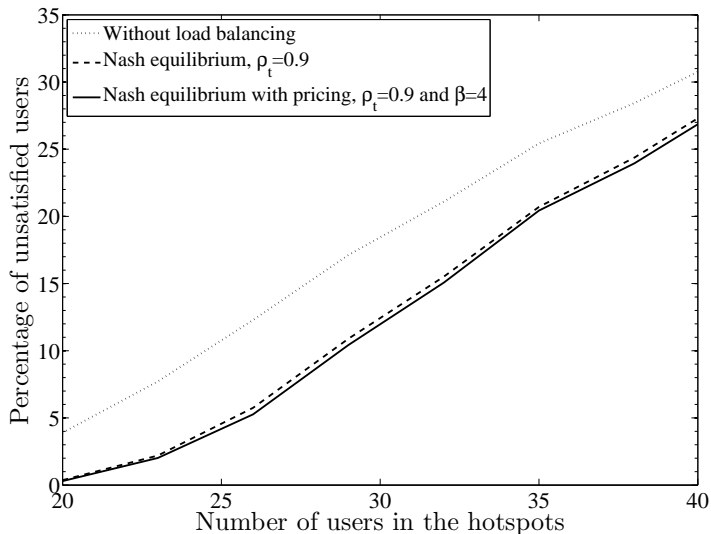
# LINEAR PRICING

- ▶ The behavior of the players can be altered by introducing a certain cost
- ▶ The players can maximize a modified utility function defined as

$$\text{utility}^c = \text{utility} - \beta \max\{N_{\text{tot}}(\rho - 1), 0\}$$

where  $\beta$  is a positive scalar that should be tuned to increase the overall performance in the network

# LINEAR PRICING



# CONCLUSION

- ▶ Non-cooperative approach in load balancing can still achieve a remarkable gain when compared to the case with no load balancing
- ▶ The overall network performance can be slightly improved using linear pricing