



CASP – A Lightweight QoS Signaling Protocol

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Lightweight QoS Signalling for Mobile Multimedia Scenarios (LQS)

- Overall project theme:
„Design of a *technology independent Lightweight QoS Signalling protocol for access networks*“
- Joint Project with members of Siemens (CT, ICN and ICM) together with Prof. Henning Schulzrinne, Columbia University
- In parallel to the joint activity in IETF WG NSIS (Next Steps In Signalling) with members from Siemens (RMR, CT, ICM), NEC, TU Berlin, Univ. of Ulm

LQS: Why are existing solutions insufficient?

- inter-domain signaling
- out-of-path signaling
- signaling other than end-to-end
- bi-directional signaling
- mobility support
- multicast support usually not needed
- interworking with policy, security, TE, ...

Siemens Requirements for LQS (I)

M – Per-flow signaling

M – Mobility support

M – Conformance to NSIS requirements / framework

=> evaluation of usability of RSVP or a descendant

=> protocol extendible to the rest of the network

M – Signaling between End System and Access Network

(M – Mandatory, O – Optional / to be discussed)

Siemens Requirements for LQS (II)

M - Support of Signaling Proxies

M - Signaling in the Access Network only

M – Heterogeneous Access Networks

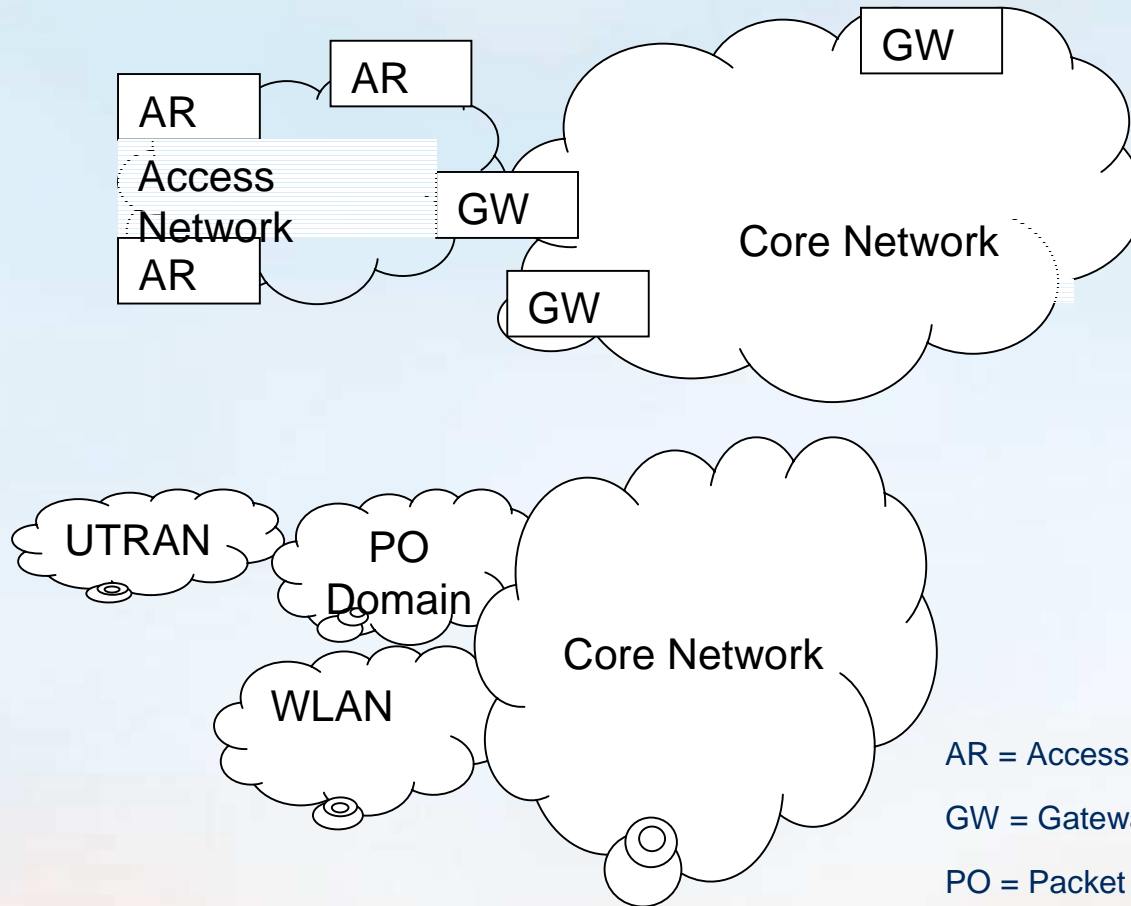
M - Independence of mobility protocols

O – Interface to other layers (e.g. support of adaptive applications, interworking with link-layer QoS)

O - Signaling across IPv4/IPv6 boundaries

(M – Mandatory, O – Optional / to be discussed)

What is an Access Network?



Core Network

- Aggregated traffic
- No mobility

Access network

- Access to end systems
- Per-flow QoS signaling
- Traffic aggregation possible
- Mobility

AR = Access Router

GW = Gateway

PO = Packet Oriented

UTRAN = UMTS Radio

Access Network

LQS: Where to use?

- QoS Signalling
- Configuration of middleboxes
- Topology discovery
- Measurement data collection
- MPLS label distribution
-

LQS: Why now?

■ Signalling solution needed:

- In 3GPP for inter-domain signaling
- Between bandwidth-brokers
- In conjunction with mobility

■ IETF NSIS (Next Steps In Signaling) WG chartered in Nov. 2001

- Generates wide interest
 - most active participants - mobility community (Siemens, NEC, Ericsson, Nokia, Alcatel)
 - closely followed and supported by „RSVP inventors“

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CASP – Cross- Application Signaling Protocol

What is CASP?

■ **Generic *signaling service***

- establishes state along path of data
- one sender, typically one receiver
 - can be multiple receivers → multicast
- *can* be used for QoS per-flow or per-class reservation
- but not restricted to that

■ **Avoid restricting users of protocol:**

- sender vs. receiver orientation
- more or less closely tied to data path
 - router-by-router
 - network (AS) path

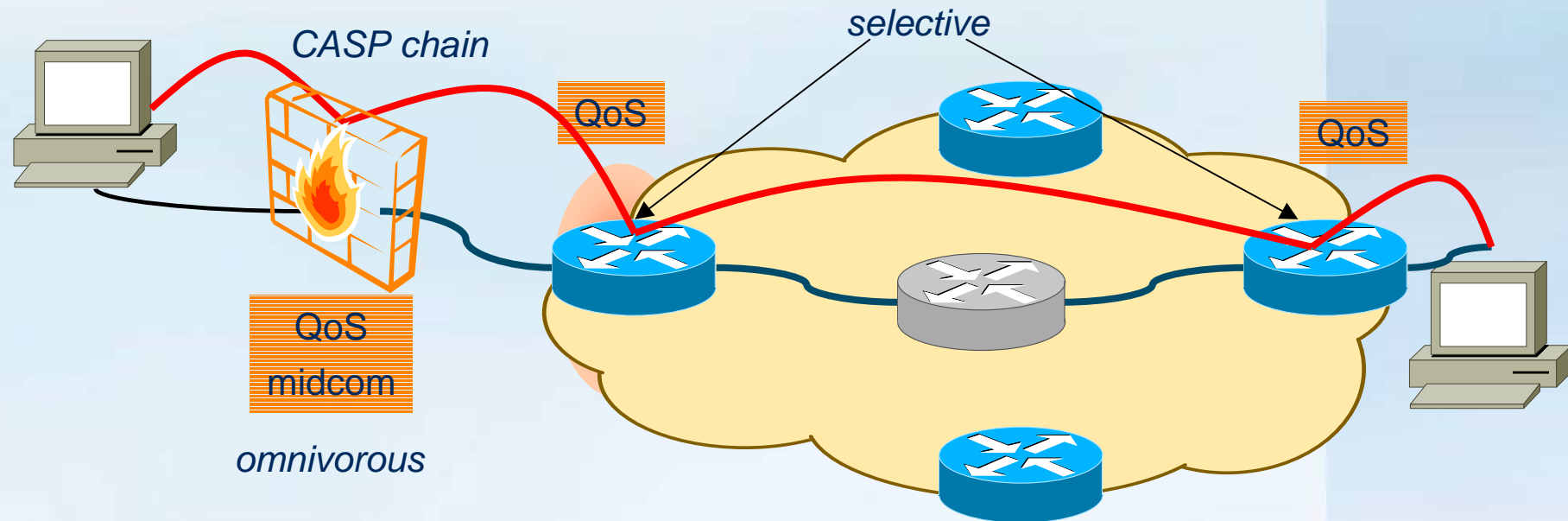
CASP Properties

- **Layered**
 - M(essaging) - layer
 - C(lient) - layer
- **Network friendly**
 - congestion-controlled
 - re-use of state across applications
- **Transport neutral**
 - any reliable protocol
 - initially, TCP and SCTP
- **Policy neutral**
 - no particular AAA policy or protocol
 - interaction with COPS, DIAMETER needs work
- **Soft state**
 - per-node time-out
 - explicit removal

CASP Properties

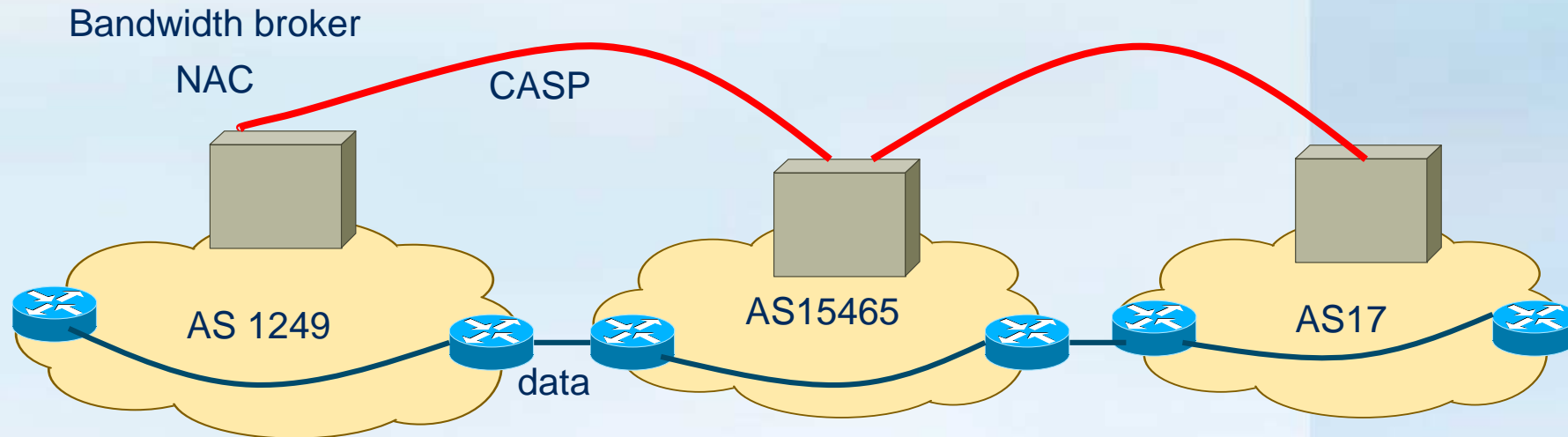
- **Extensible**
 - data format
 - feature negotiation
- **Security protection**
 - first peer, intra and inter domain
- **Topology hiding**
 - hide addresses of visited nodes
- **Light-weight**
 - message forwarding overhead
 - implementation complexity
- **Mobility transparent**

CASP Network Model – On-Path



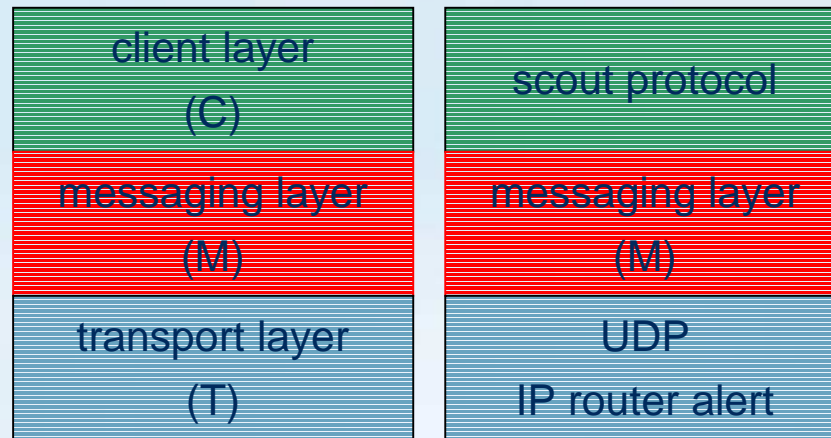
- **CASP nodes form CASP chain**
- **Not every node processes all client protocols:**
 - non-CASP node: regular router
 - omnivorous: processes all CASP messages
 - selective: bypassed by CASP messages with unknown client protocols

CASP Network Model – Out-of-Path



- Also route network-by-network
- Can combine router-by-router with out-of-path messaging

CASP Protocol Structure



■ Client layer does the real work:

- reserve resources
- open firewall ports
- ...

■ Messaging layer:

- establishes and tears down state
- negotiates features and capabilities

■ Transport layer:

- reliable transport
e.g. TCP, SCTP

CASP Messages

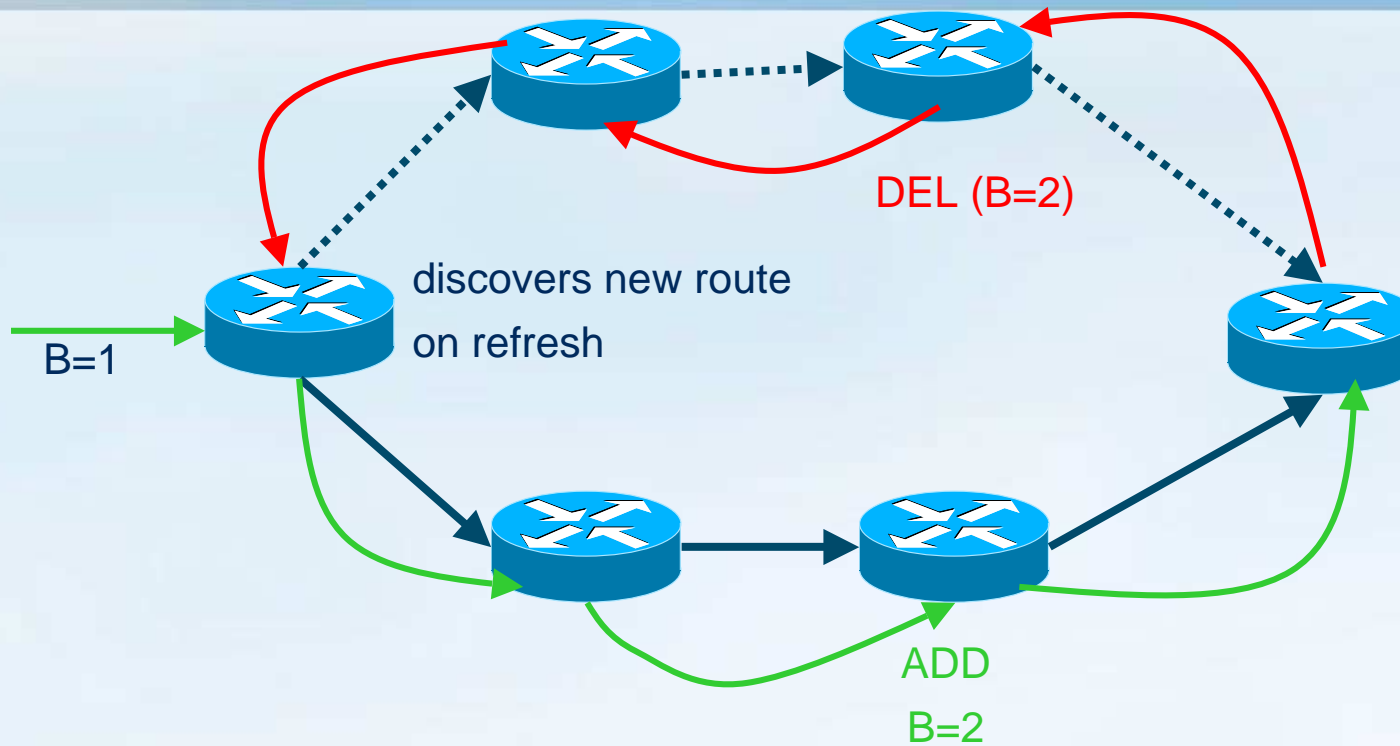
- **Regular CASP messages**
 - establish or tear down state
 - carry client protocol

- **Scout messages**
 - discover next hop

- **Hop-by-hop reliability**

- **Generated by any node along the chain**

CASP: Mobility and Route Changes



- Avoids session identification by end point addresses
- Avoid use of traffic selector as session identifier
- Remove dead branch

- **Security for the M(essaging) layer**
 - IPSec
 - TLS (possibly with EAP on top of it)
 - many different key exchange protocols supported (IKE, KINK, SOI, etc.)

- **Security for the C(lient) layer**
 - based on the security of the M-layer
 - CMS used to selectively wrap objects and to provide protection for them

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CASP QoS Client Protocol

CASP QoS Client Protocol Feature Overview

- **Direction Neutral**
- **Bidirectional Reservation**
- **Reservation Range**
- **Partial Reservation**
- **Advance Reservation**
- **Resource Query**
- **Reserve / Commit Mechanism**
- **Local Information**

CASP QoS Client Protocol Features

- **Direction Neutral**
 - sender-oriented and receiver-oriented reservations
- ✦ **Message reduction**
- ✦ **Support for adaptive application**

- **Bidirectional Reservation**
 - single reservation for symmetric routes
 - support of asymmetric reservation
- ✦ **Message reduction**
- ✦ **Enhanced resource utilization**

CASP QoS Client Protocol Features

- **Reservation Range**
 - specification of upper and lower resource threshold
- ✦ **Message reduction**
- ✦ **Support for adaptive application**

- **Partial Reservation**
 - application may accept „blackholes“ for some time
- ✦ **Enhanced resource utilization**

- **Advance Reservation**
 - resource reservation at any time in future
- ✦ **Potential feature for conference applications**

CASP QoS Client Protocol Operation

- **Resource Query / Response**
 - query resources before requesting them
- ★ **Avoid unnecessary resource allocation**

- **Reserve / Commit Resources**
 - reserve assigns resources
 - commit allocates resources for exclusive use
 - receiver / Sender oriented reservation
 - priority object for reservation priority
- ★ **Efficient resource usage**

- **Local information**
 - authentication, DSCP, accounting
- ★ **Minimize information overhead**

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**CASP – Future
Work
Plans and Ideas**

CASP – Future work

■ Implementation

- Message layer (in-band, hop-by-hop)
 - likely, Linux or FreeBSD
- Discovery mechanisms
 - routing-based (OSPF)
 - scout protocol
- QoS client
- out-of-band messaging
 - "bandwidth broker" or NAC model
- integrate with traffic control

■ Enhancements

- investigate tunnels

CASP future work

- **Performance analysis**
 - message handling
 - TCP and TLS set-up overhead
 - maximum number of simultaneous connections
- **Specify additional client protocols**
 - NAT and firewall control ("midcom")
 - MPLS or lightpath setup?
 - denial-of-service traffic filter?
 - router QoS management → gather performance statistics
- **Interaction with AAA**
 - authentication, authorization and accounting
 - something other than COPS?
- **IETF**
 - submit Internet Draft to IETF NSIS working group
 - pursue standardization